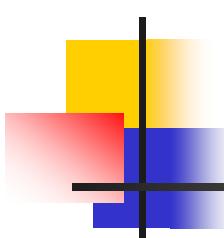


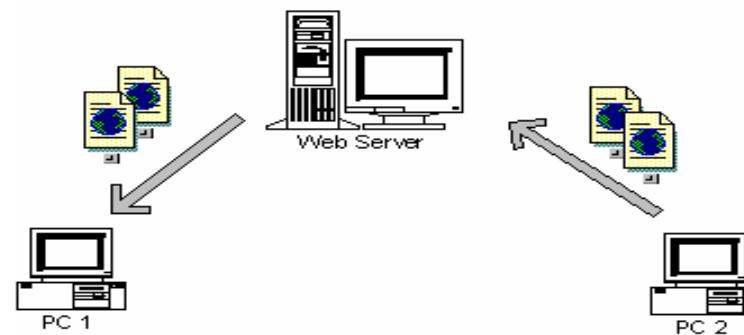
# Data Management in P2P Systems: Challenges and Research Issues

Timos Sellis

National Technical Univ. of Athens



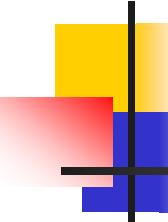
# What is P2P?



Client Server Architecture



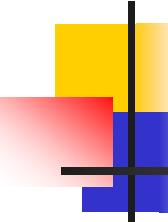
Peer-to-Peer Architecture



# P2P: an old concept

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- P2P is not a new technology
  - *Oct. 29 1969: first transmission UCLA → Stanford Research Inst. (SRI) [UCSB, U. Of Utah]; Peer computing status among independent computing sites*
- 1970s-1980s: ARPANET completed (1978); Usenet and DNS
- 1990s: NSFnet and Internet explosion
- 1996: ICQ (bypassing DNS by offering own addressing scheme and allowing end points to be directly addressable with each other)

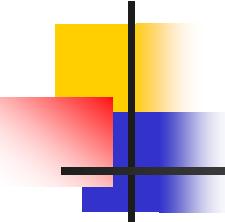


# NAPSTER boosts P2P technology

- 1999: Napster is born
  - universities notice first
  - mid-2001: shut down service; other apps emerge since 2000
- 2002: Napster is gone...
- April 30, 2002:

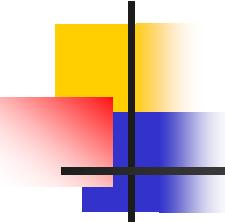
A study from WebSense indicates that the number of file-swapping, and peer-to-peer websites, has grown by 535 percent in the past year, despite legal efforts to have them shut down.

According to the study findings, the number of p2p websites totals nearly 38,000.



# The Evolution of P2P systems

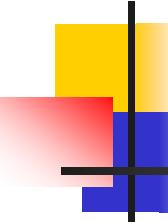
- First generation – centralized P2P systems
  - *E.g. Napster, SETI@home*
- Second generation –decentralized & unstructured P2P systems
  - *E.g. Gnutella*
- Third generation—structured P2P systems
  - *DHT systems (CAN/Chord/Pastry/Tapestry)*
  - *Skip-list based systems*
- ....



# P2P Perspective: Self-Organized Systems

P2P Paradigm is characterized by:

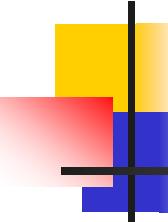
- no centralized control
- adaptable behavior
- self-organization
  - *distribution of control  $\Rightarrow$  decentralization*
  - *local interactions, information and decision  $\Rightarrow$  autonomy*
- emergence of global structures
- failure resilience and scalability



# P2P Paradigm Dimensions

---

- Structure – Framework
  - *Servers and clients*
- Organization
  - *Structured vs. unstructured overlays*
- Managed and stored data
  - *Files vs. (semi)-structured data*
- Functionality-Applications
  - *Sharing, resource management, web services, collaboration etc*



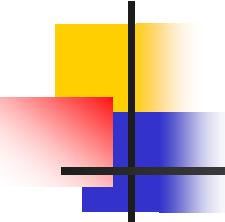
# Structure- Frameworks

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Main differences of different structures:

*how information is shared and how much information is shared*

- Models:
  - *Centralized*
  - *Decentralized*
  - *Controlled-decentralized | Hybrids*
- Components:
  - client, server, servents



## Centralized: pros and cons

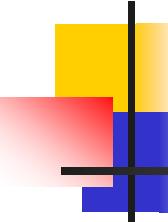
Centralized servers keep a catalog of available files ⇒ Napster

### Pros

- More effective, comprehensive searches
- Access is controlled

### Cons

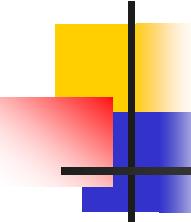
- System has single points of entry; one fail could bring the whole system down
- Broken links, out of date information.



# Decentralized Framework

---

- Every user acts as a client, a server or both (servent).
- User connects to framework and becomes a member of the community, allowing others to connect through him/her
- Users “talk” directly to other users with no intermediate or central authority
- No one entity controls the information that passes through the community  
    ⇒Gnutella, Freenet



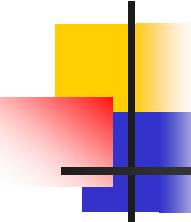
# Decentralized: Pros and Cons

## Pros

- Peers “talk” directly with no central authority.  
Nobody owns the Gnutella Network and nobody can shut it down.
- Isolated node failure can quickly and automatically be worked around.
- Free loading
- Scalability

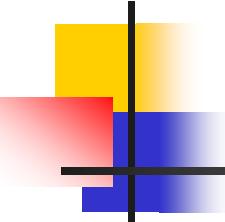
## Cons

- Searches are less effective and can be slow.



# Hybrids: Controlled decentralization

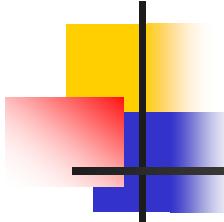
- Characteristics of both centralized and decentralized frameworks, i.e. systems with *super-peers*
- User's computer may act as **client**, a **server** or a **servent**; there are server operators which may control which clients and/or servents are allowed to access a particular server.
- Examples: Morpheus, Groove



# Organization: Overlay Networks

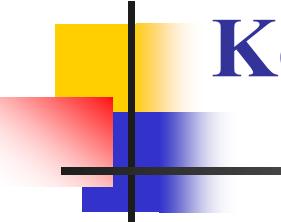
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- Structured vs. unstructured overlays
- Structured:
  - data is indexed and stored at peers with pre-specified characteristics, e.g. peer id
  - Peers are connected to specific peers
- Unstructured:
  - No pre-specified connection of peers and storage/indexing of data



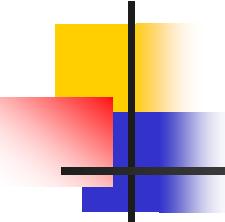
# Managed and Stored Data

- Simple files:
  - search is performed with keywords for file names and file extensions
- Structured or semi-structured data ⇒ *P2P Databases*:
  - XML data
  - Relational data
  - RDF data



# Key Application/Function Areas

- File sharing / Content distribution:
  - *Not just mp3s/media sharing*
  - *Distributed searching: Used to easily lookup and share files and offer content management*
- Instant messaging:
  - *Exchanging messages and chatting ⇒ Jabber, IM*
- Distributed computation:
  - *Use under utilized Internet and/or network resources for improving computation and data analysis ⇒ Grid Computing*
- P2P groupware / Collaboration:
  - *Cooperative publishing, messaging, group project management. Secure environments are offered in some products*

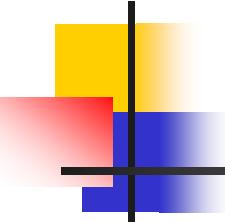


# What Has to Be Done

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Issues to be worked out:

- Reliability  
*Is the provided content or the provider itself reliable? How can we assure these?*
- Resource management  
*How can we achieve efficient collaboration of autonomous resources?*
- Security & Privacy  
*How can we guarantee security together with privacy without rescinding peer autonomy?*
- Intellectual Rights  
*Legal notions of copyright and intellectual property need to adapt to the sharing model of the P2P paradigm*



# Research Areas

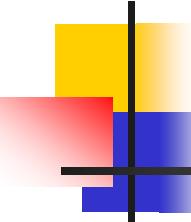
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- Peer discovery and group management
- Data location and placement
- Reliable and efficient file exchange
- Security/privacy/anonymity/trust



# Relational data sharing in Unstructured P2P vs. Distributed DB

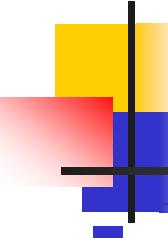
P2P	Distributed Database Systems
Nodes can join and leave the network anytime.	Nodes are added/removed from the network in a controlled manner.
Usually no predetermined (global) schema among nodes. Queries: Keywords	Have some knowledge of a shared schema. Queries: SQL
Answers to queries are typically incomplete (by “completeness” we mean all answers that satisfy a query)	Can actually retrieve the complete set of answers.
Content location is typically by “word-of-mouth” e.g., node routes query to its neighbors and so on...	Exact location to direct the query is typically known.



# P2P & DB Systems

P2P      DB

Flexibility	✓	✗
Decentralized	✓	✗
Strong Semantics	✗	✓
Powerful query facilities	✗	✓
Fault Tolerance	✓	✗
Lightweight	✓	✗
Transactions & Concurrency Control	✗	✓

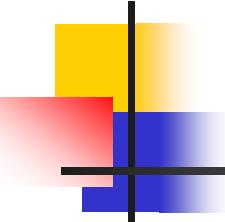


# P2P + DB = ?

## ■ P2P Database? No!

- *ACID transactional guarantees do not scale, nor does the everyday user want ACID semantics*
- *Much too heavyweight of a solution for the everyday user*
- **Query Processing on P2P!**
  - *Both P2P and DBs do data location and movement*
  - *Can be naturally unified (lessons in both directions)*
  - *P2P brings scalability & flexibility*
  - *DB brings relational model & query facilities*

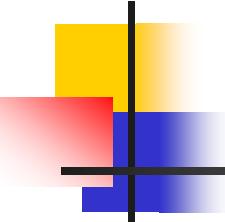
Taken from Hellerstein's group ppt



# Many New Challenges

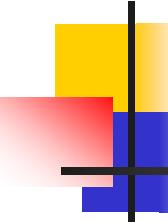
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- Relative to other parallel/distributed systems
  - *Partial failure*
  - *Churn*
  - *Few guarantees on transport, storage, etc.*
  - *Huge optimization space*
  - *Network bottlenecks & other resource constraints*
  - *No administrative organizations*
  - *Trust issues: security, privacy, incentives*
- Relative to IP networking
  - *Much higher function, more flexible*
  - *Much less controllable/predictable*



# Data management in P2P systems

- Current research focuses on
  - *Decentralized schema mappings*
    - PeerDB: unstruct. network, keyword search only
  - *Extending DHT for complex querying*
    - PIER : exact-match and join queries
  - ***Query reformulation***
    - Edutella: super-peer, RDF-based schemas
    - Piazza: graph of pair-wise schema mappings
  - *Replication*
    - generally limited to static read-only files
    - P-Grid addresses updates in structured networks



## Our model of P2P databases

---

- All peers hold relational databases
- There is no mediation or “global schema”
- Peers only have a few **acquaintees** (neighbors), i.e. other peers they “know” and can “talk to”
- Peers receive queries, answer them and at the same time **propagate** them to their neighbors
- Peers use mappings to map their own schema to their neighbor’s schemas in order to correctly translate a query and propagate it (Query rewriting)
- There is ad-hoc and intermittent connectivity

# Query Answering: Motivating Application

Patients(pid, name, age, sex)

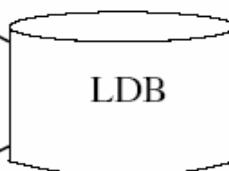
Treatments(pid, date, symptoms, diseasedescr, treatrec, drugname)



```
SELECT  
age,sex,symptoms,treatrec,drugname  
FROM Patients,Treatments  
WHERE Patients.pid = Treatments.pid  
and Treatments.diseasedescr = "X"
```

$Q_{orig}$

```
SELECT age,sex  
FROM Patients,Results  
WHERE Patients.pid = Results.pid and  
Results.diagnosis = "X"
```



Patients(pid, name, age, sex)  
Results(testid, pid, date, results, diagnosis)

$Q_{orig\_LDB}$

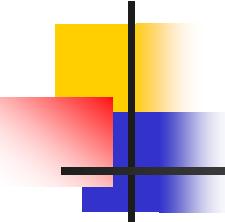


```
SELECT age,sex  
FROM Patients,Treatments  
WHERE Patients.pid = Treatments.pid and  
Treatments.diseasedescr = "X"
```

$Q_{LDB\_StuartDB}$

Patients(pid, name, age, sex)

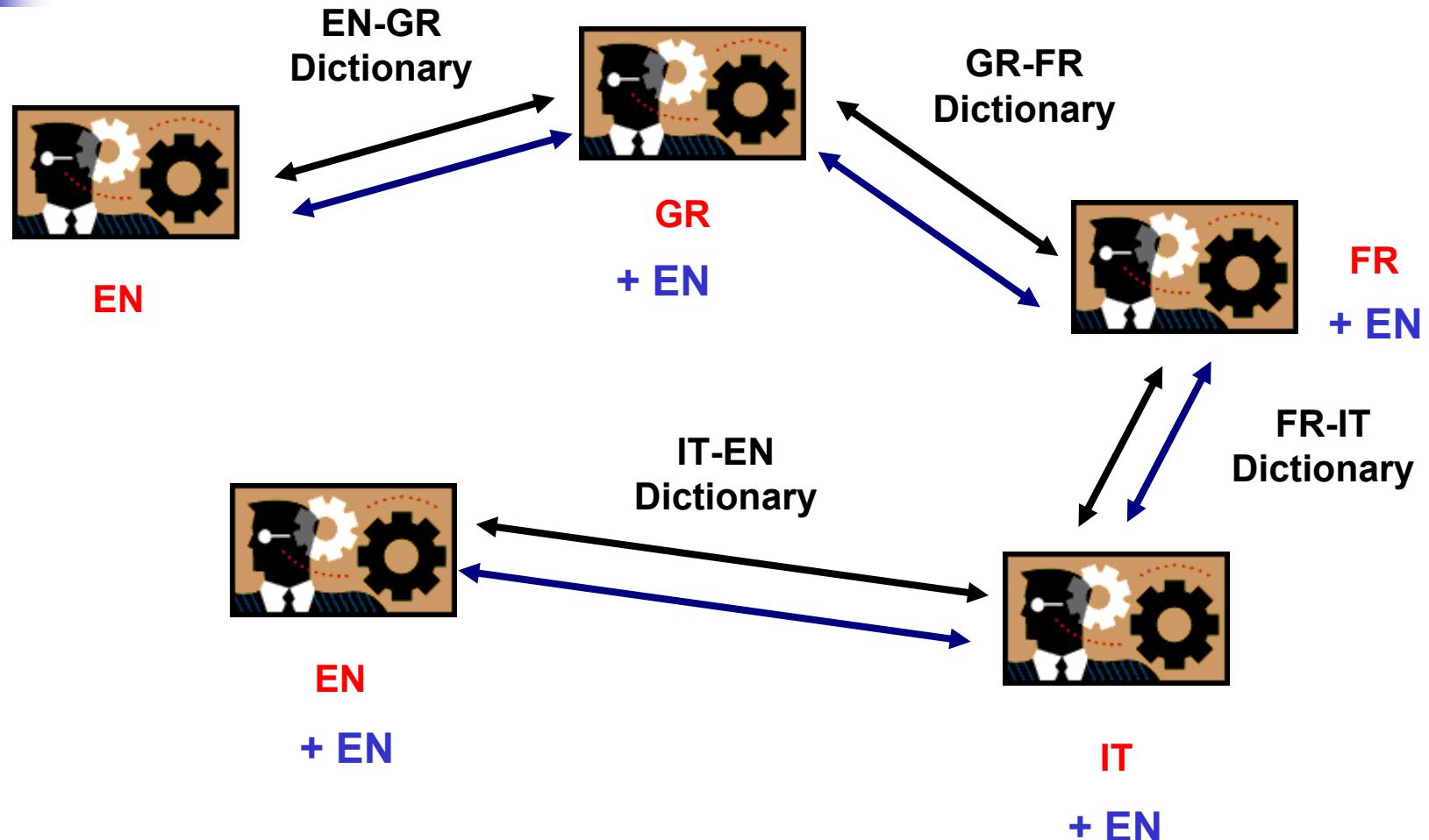
Treatments(pid, date, symptoms, diseasedescr, treatrec, drugname)

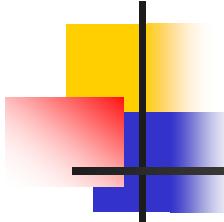


## Traditional Query Propagation (Contd.)

- Each peer that receives a query  $Q$ 
  - *Rewrites it according to its local schema ( $Q'$ )*
  - *Answers  $Q'$*
  - *Forwards  $Q'$  to one or more of its neighbors*
- Peers with dissimilar schemas hide nodes rich in information
- Bad rewritings have the same effect
- Local-only knowledge in the overlay puts constraints on the propagation of a query
  - *Bandwidth limitations*
  - *Cannot keep track of all nodes!*

# People have a similar problem





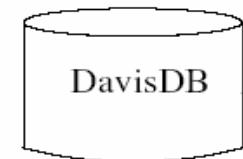
# Overview of Our Framework - GrouPeer

- Forward **both** the **original** and the **rewritten** version of Q
- Answer the one most similar to the local schema
  - *Similarity function*at the same time .....
- Overlay restructuring according to evaluation of the results (learning new neighbors)  
but do this in an.....
- Intelligent, bandwidth-efficient query forwarding with informed walks

# Answering the Original Query

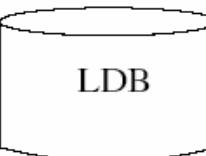
Patients(pid, name, age, sex)

Treatments(pid, date, symptoms, diseasedescr, treatrec, drugname)



```
SELECT  
age,sex,symptoms,treatrec,drugname  
FROM Patients,Treatments  
WHERE Patients.pid = Treatments.pid  
and Treatments.diseasedescr = "X"
```

$Q_{\text{orig}}$



Patients(pid, name, age, sex)  
Results(testid, pid, date, results, diagnosis)



$Q_{\text{LDB\_StuartDB}}$

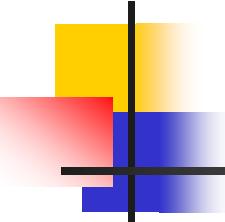
~~SELECT age,sex  
FROM Patients,Treatments  
WHERE Patients.pid = Treatments.pid and  
Treatments.diseasedescr = "X"~~

$Q_{\text{orig\_StuartDB}}$

~~SELECT  
age,sex,symptoms,treatrec,drugname  
FROM Patients,Treatments  
WHERE Patients.pid = Treatments.pid and  
Treatments.diseasedescr = "X"~~

Patients(pid, name, age, sex)

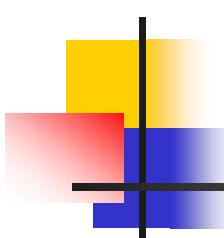
Treatments(pid, date, symptoms, diseasedescr, treatrec, drugname)



# Similarity Measure

---

- Simple structural similarity between 2 versions of Q:  
 $M_{sim}(Q_{v1}, Q_{v2})$  correlates simple conceptual similarities of structural query elements (i.e. ‘select’ attributes and ‘where’ conditions )
  
- Different at each peer
  - *Schema matching methodology*
  - *Individual preferences*



# Confidence Parameter $\theta$

---

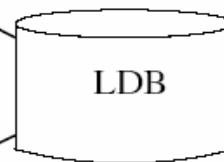
- Express each peer's individuality
  - *Schema matching ability*
  - *Query structure*
  - $\theta = \theta_p(Q)$
- Each peer decides which query to answer (original, rewritten) according to:

$$Cov_p(Q_2, Q_1) = M_{sim\ p}(Q_1, Q_2) \cdot \theta_p(Q_1)$$

# Example Query Answering

Patients(pid, name, age, sex)

Treatments(pid, date, symptoms, diseasedescr, treatrec, drugname)



Patients(pid, name, age, sex)  
Results(testid, pid, date, results, diagnosis)

$$\text{Cov}_{\text{StuartDB}}(Q_{QLDB\_Stuart}, Q_{\text{orig}}) = 0.4 \cdot 1 = 0.4$$



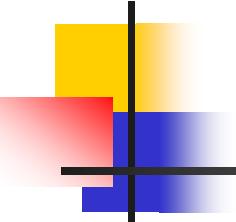
SELECT age,sex  
FROM Patients,Treatments  
WHERE Patients.pid = Treatments.pid and  
Treatments.diseasedescr = "X"

Patients(pid, name, age, sex)

Treatments(pid, date, symptoms, diseasedescr, treatrec, drugname)

$$\text{Cov}_{\text{StuartDB}}(Q_{\text{orig\_Stuart}}, Q_{\text{orig}}) = 1.0 \cdot 0.5 = 0.5$$

SELECT  
age,sex,symptoms,treatrec,drugname  
FROM Patients,Treatments  
WHERE Patients.pid = Treatments.pid and  
Treatments.diseasedescr = "X"



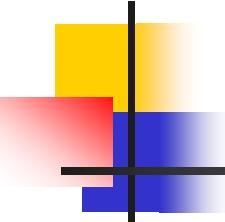
# Learning Process

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- Data providers return to the query originator:
  - **Which** version they answered
  - **Successively rewritten** query
  - **Locally rewritten** query
  - **Answer** tuples
- The originators evaluate the answers:

$$Ev(Ans) = Cov(Q_x, Q) \cdot \max(1, w_Q \cdot |Res(Q_x)|)$$

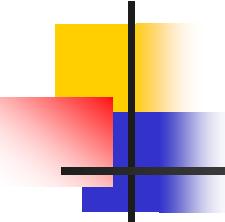
$$Cov(Q_x, Q) = M_{sim}(Q, Q_x) \cdot \theta_{local}(Q)$$



## Semantic Grouping of Peers

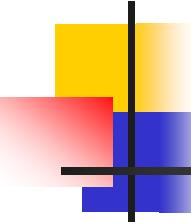
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- The learning process leads to semantic clustering of the P2P network
- Each semantic cluster can create an explicit interest group:
  - *Group Initiator: Peer that initiates the group creation*
  - *Group Schema: An abstract schema with most common schema elements of the cluster*



# Routing and Restructuring in the Overlay

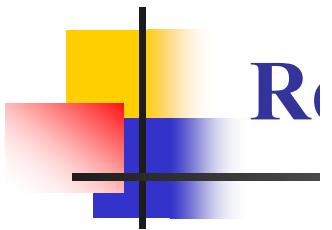
- Informed walks instead of blind forwarding
  - $K$  walkers, TTL steps each
  - Choose neighbors according to schema similarity
- Adding an acquaintance p:
  - When  $\theta_p = \theta_{local}$
  - More than  $T$  replies from p
  - Impose a maximum number of neighbors
- Dropping an acquaintance
  - Low(est) similarity neighbors
  - Only if they don't get disconnected



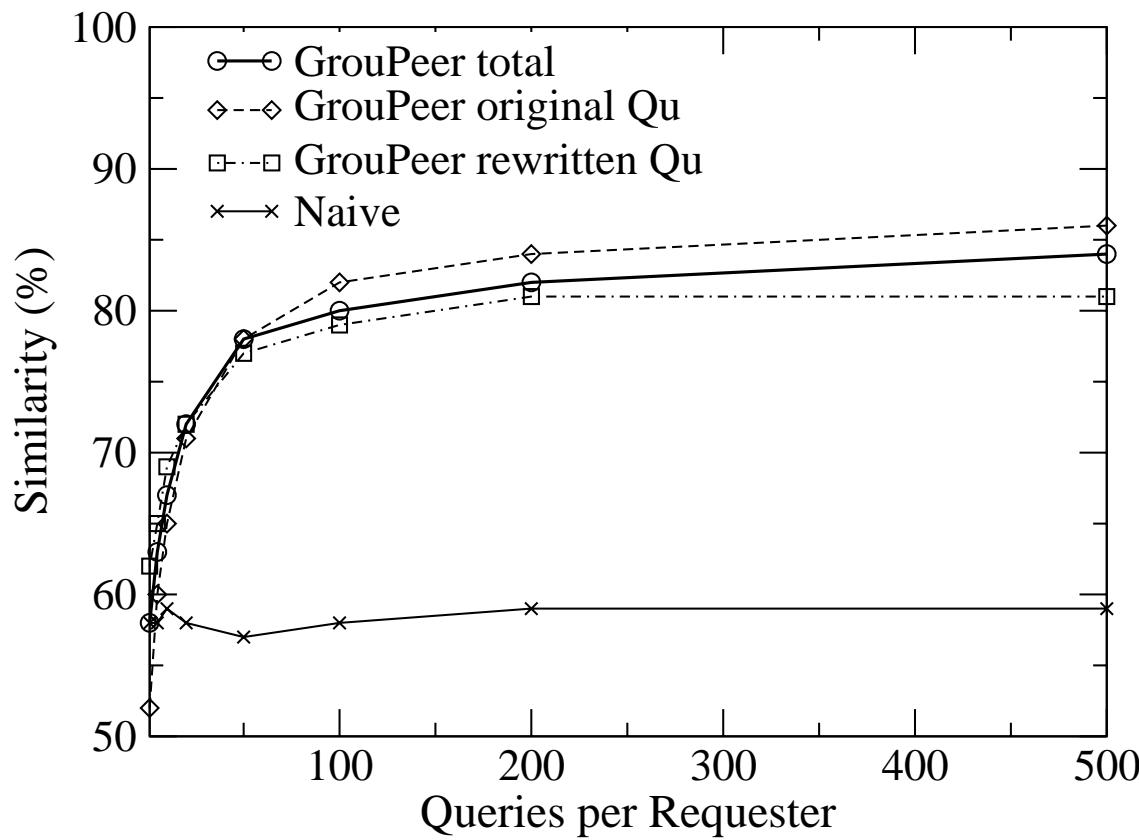
# Simulation Methodology

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- Simulator in C
- Pure P2P model, 1-10K graphs (power-law and random)
- 100 attributes, random assignment (10 avg.)
- Queries on local schemas (3 attributes avg.)
- 100 requesters
- Compared methods
  - *Naïve rewriting (naïve)*
  - *Flooding with each peer answering  $Q_{orig}$  (Brute) – and some variations*

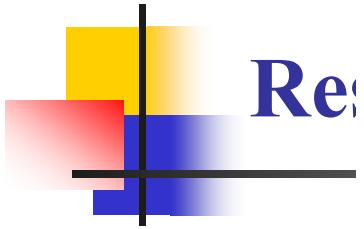


Re

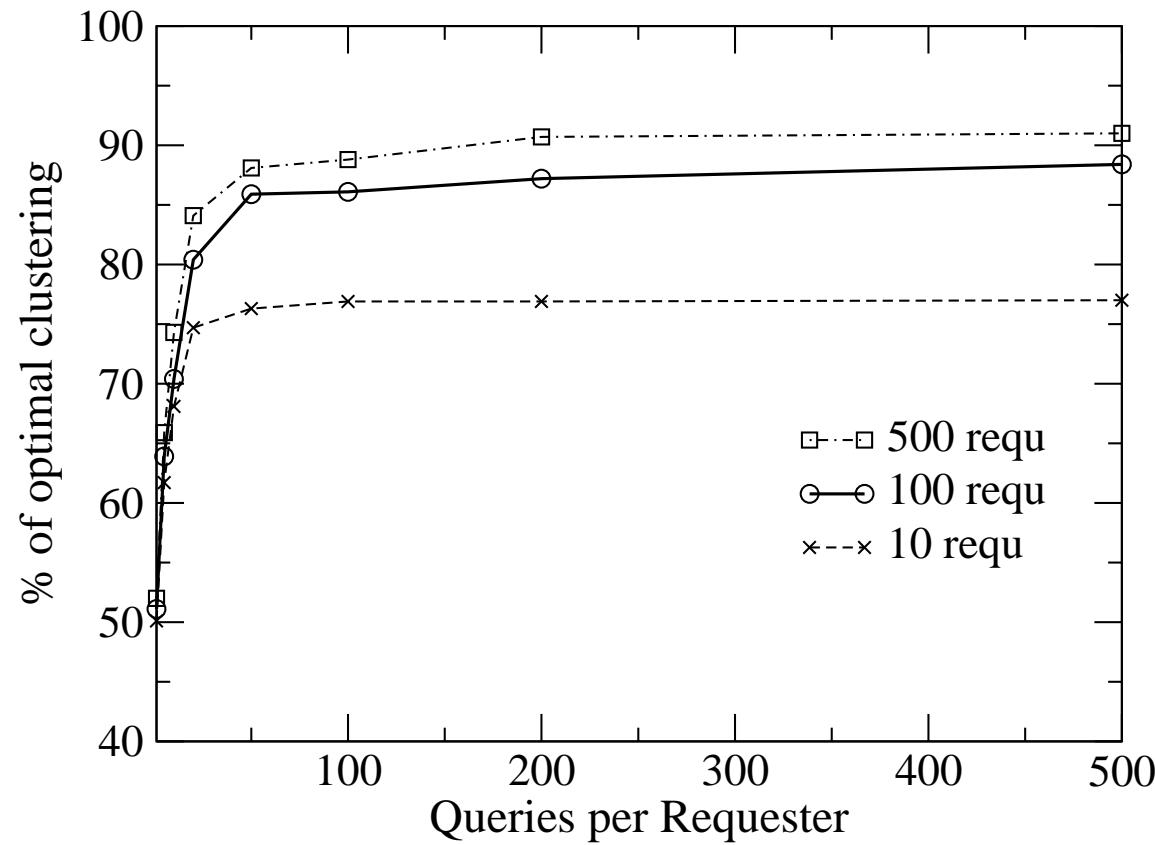


Similarity of results to the original query over variable number of queries per requester

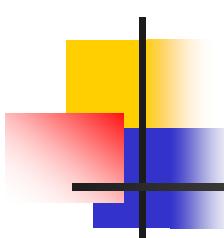
Timos Sellis-I



Re



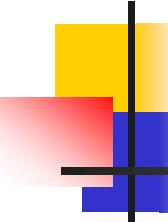
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## Related Work

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- PeerDB
- Query Reformulation (Halevy et al)
- Edutella
- Chatty Web

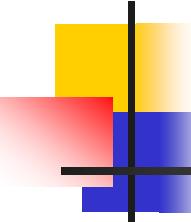


# Handling Spatial Data

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## Problem Definition

- Assumptions:
  - *Dissemination of spatial data in autonomous nodes*
  - *Structured flat P2P overlay*
  
- Problem:
  - *Necessity for indexing and routing techniques for such an environment*

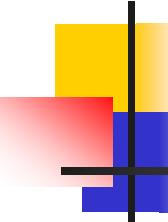


# Technique Requirements

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We need a technique that:

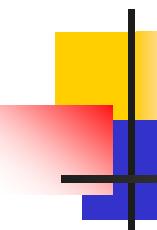
- Guarantees the retrieval of any existing spatial information in the system
- Achieves satisfying index size for any node
- Achieves a satisfying length for any search path
- Provides easier access to popular information
- Preserves proximity of areas



## Related Work

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- VBI-tree (ICDE '05)
- P2PR-tree (Workshop of EDBT '04)
- Quad-tree approach (Poster in ICDE'05)
- kd-tree approach (WebDB'04)

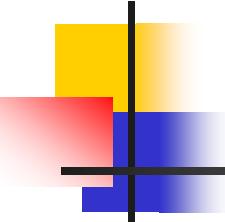


# Specification of Spatial Areas

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## Partitioning Space:

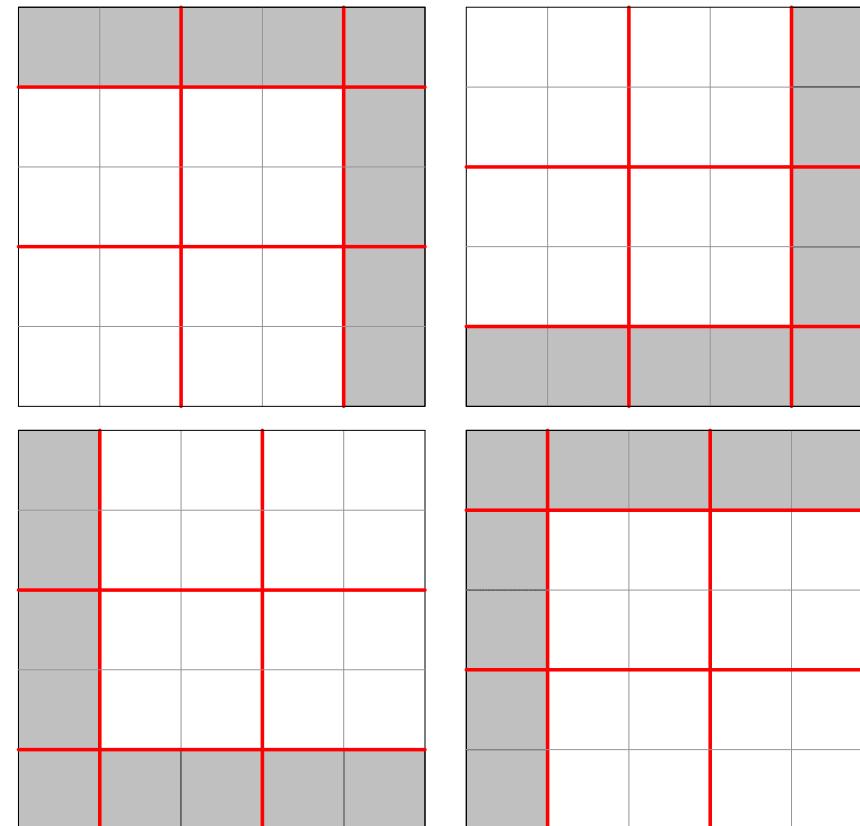
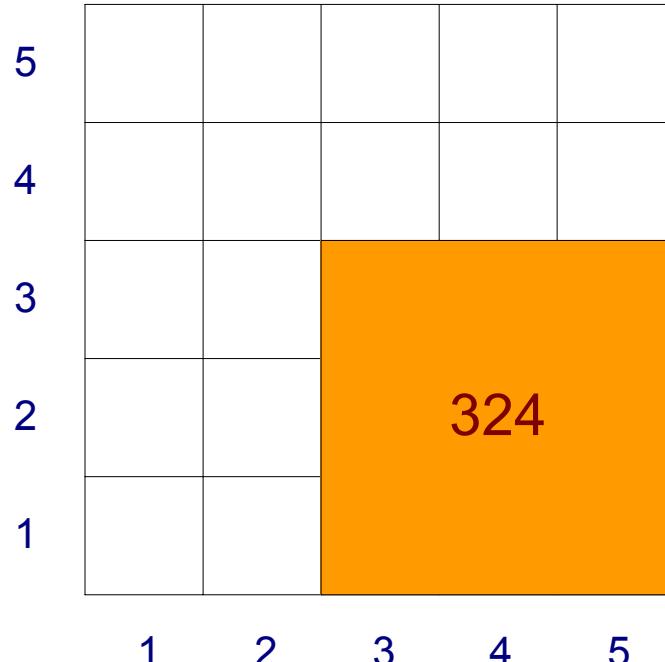
- We assume a basic fine-grained  $n \times n$  grid
- Specifying Spatial Areas:
- We consider any  $i \times i$  area on the basic grid,  $1 \leq i \leq n$   
     $\Rightarrow$  Each  $i \times i$  area belongs to an  $i \times i$  grid
- The id of an area A is:  $A \equiv a_z a_x a_y$
- $a_z$ : is the granularity level, i.e. the size in terms of basic cells.
- $a_x$ : x coordinator of the center basic cell.
- $a_y$ : y coordinator of the center basic cell.

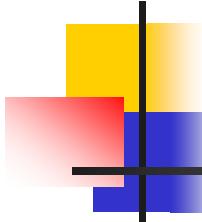


# Grid & Area Examples

The four 2<sup>nd</sup> level grids ( $z=2$ )

Basic 5x5 grid





# Distributed Indexing for Spatial Areas

Using Chord as the basic DHT we index areas as follows:

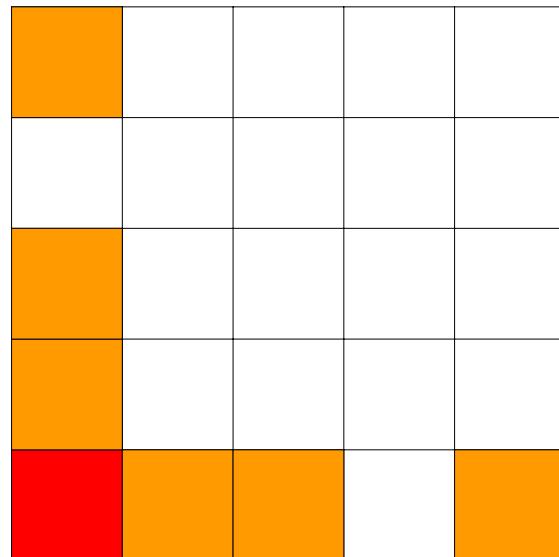
## ***Distributed Indexing Algorithm***

- **Step 1**: Indexing the grids of each level
  - *Phase A: Indexing areas of the same grid*
  - *Phase B: Indexing areas of other grids of the same level*
- **Step 2**: Indexing the grids of other levels

# Indexing Examples

Example of indexing on the basic grid with Chord of base 2

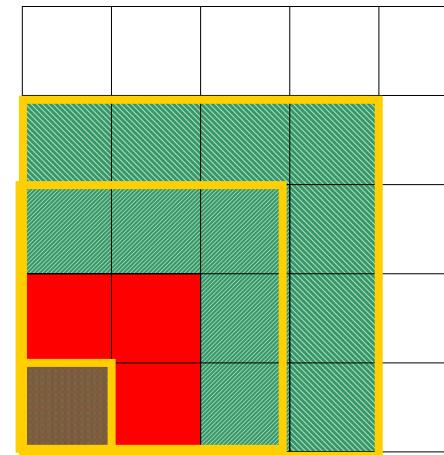
Basic 5x5 grid



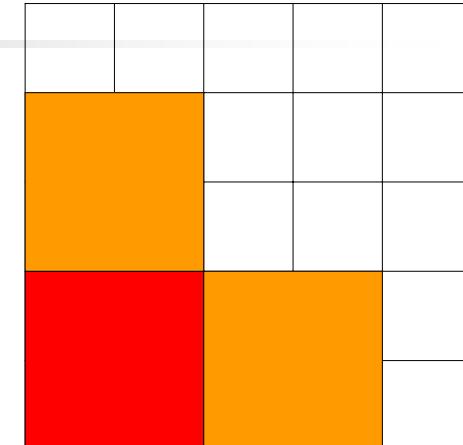
Area stored in the indexing peer

Example of indexing on a 2<sup>nd</sup> level grid with Chord of base 2

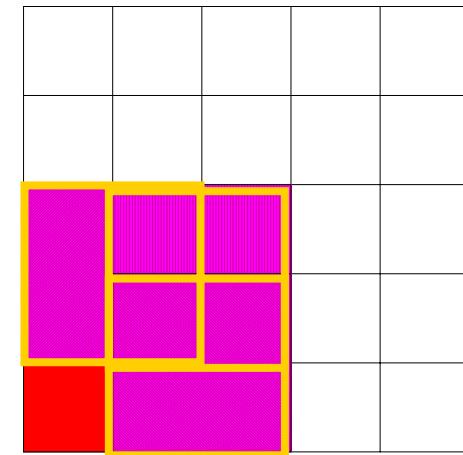
Indexing areas of different size

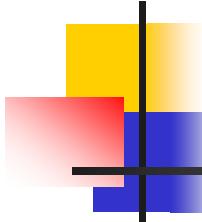


Indexing non-overlapping areas of the same size



Indexing overlapping areas of the same size





# Direct Neighbors & Data Hashing

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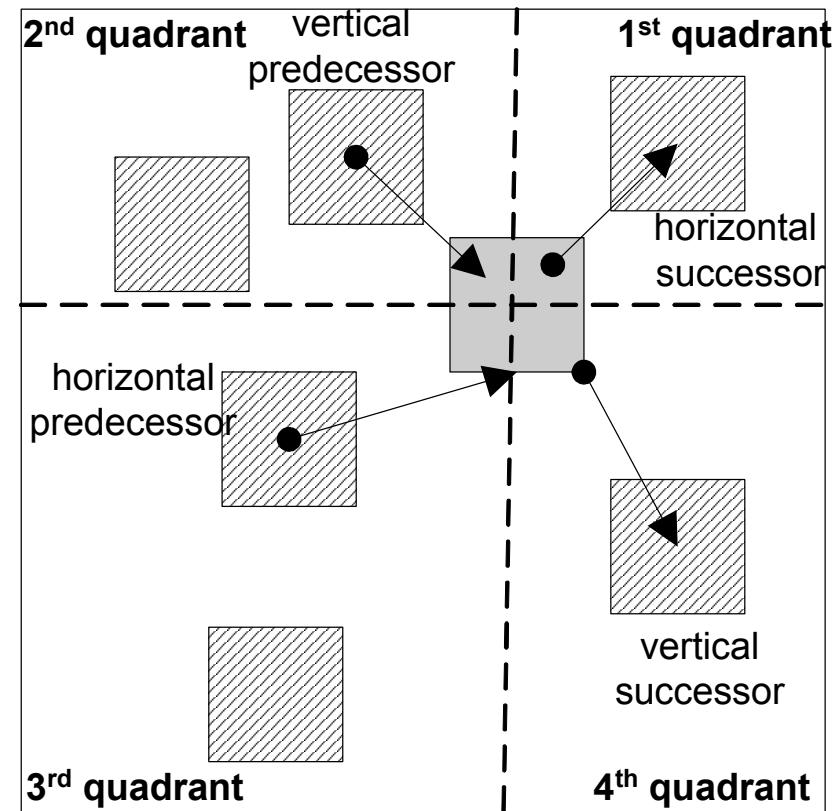
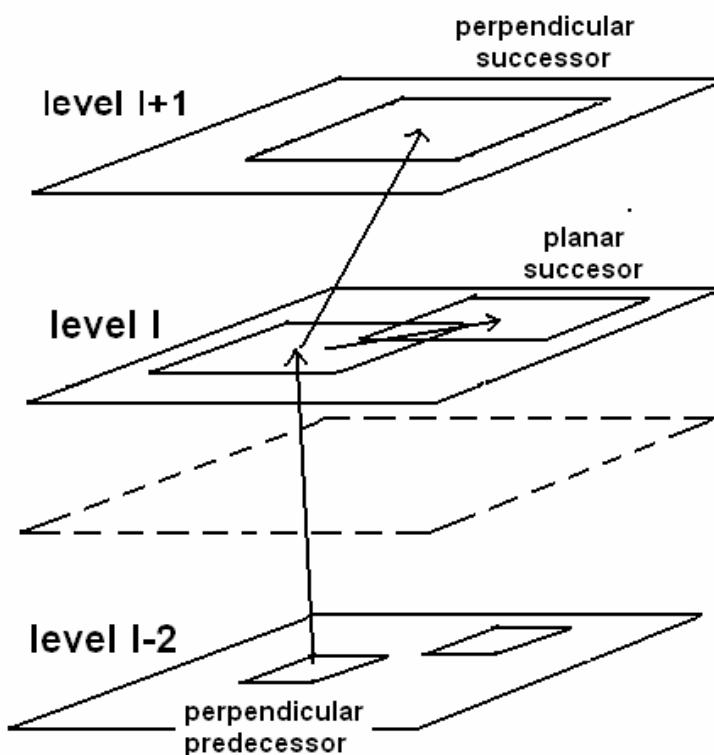
Peers have neighbors on each dimension:

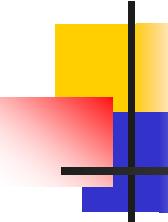
- *Horizontal*
- *Vertical*
- *Perpendicular*

*Neighbors are the closest peers on each direction.*

*Areas are stored to the closest neighbor following a priorities of dimensions*

# Direct Neighbors: Examples





# Routing of Spatial Queries

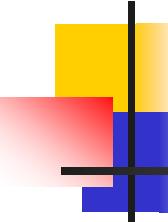
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## *Basic Routing Algorithm*

**Step 1:** Find a peer corresponding to the same size as the sought area.

**Step 2:** Find a peer hosting the sought area or an overlapping area of the latter

**Step 3:** Find the peer hosting the sought area

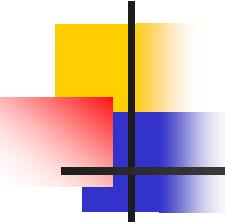


# Complexity Characteristics

Complexity for both pure Chord and our method is  $O(3\log_2 n)$ .

**But with our method:**

- Most Important: *relatively* close areas are stored in close peers
- The search path grows with the *relative* distance
- Peers with big areas are favored with small indexes
- Shorter search paths for bigger areas



## Summary

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- P2P systems offer effective use of distributed resources
- Standards for P2P infrastructure are emerging
  - *Sun (JXTA) and Microsoft (.NET)*
- Content choice and control
  - *The user becomes not only a consumer but a content provider; P2P allows the end user to participate in the Internet, as the latter was originally envisioned*

## AN EXCITING NEW AREA FOR DATABASE MANAGEMENT