

GBKR

A **G**raph-**B**ased **K**nowledge **R**epresentation
and Reasoning Model

Agenda

- Knowledge Representation and Reasoning
- GBKR model
- Relationships with Logics
and other computational models

Knowledge Representation & Reasoning ⁽¹⁾

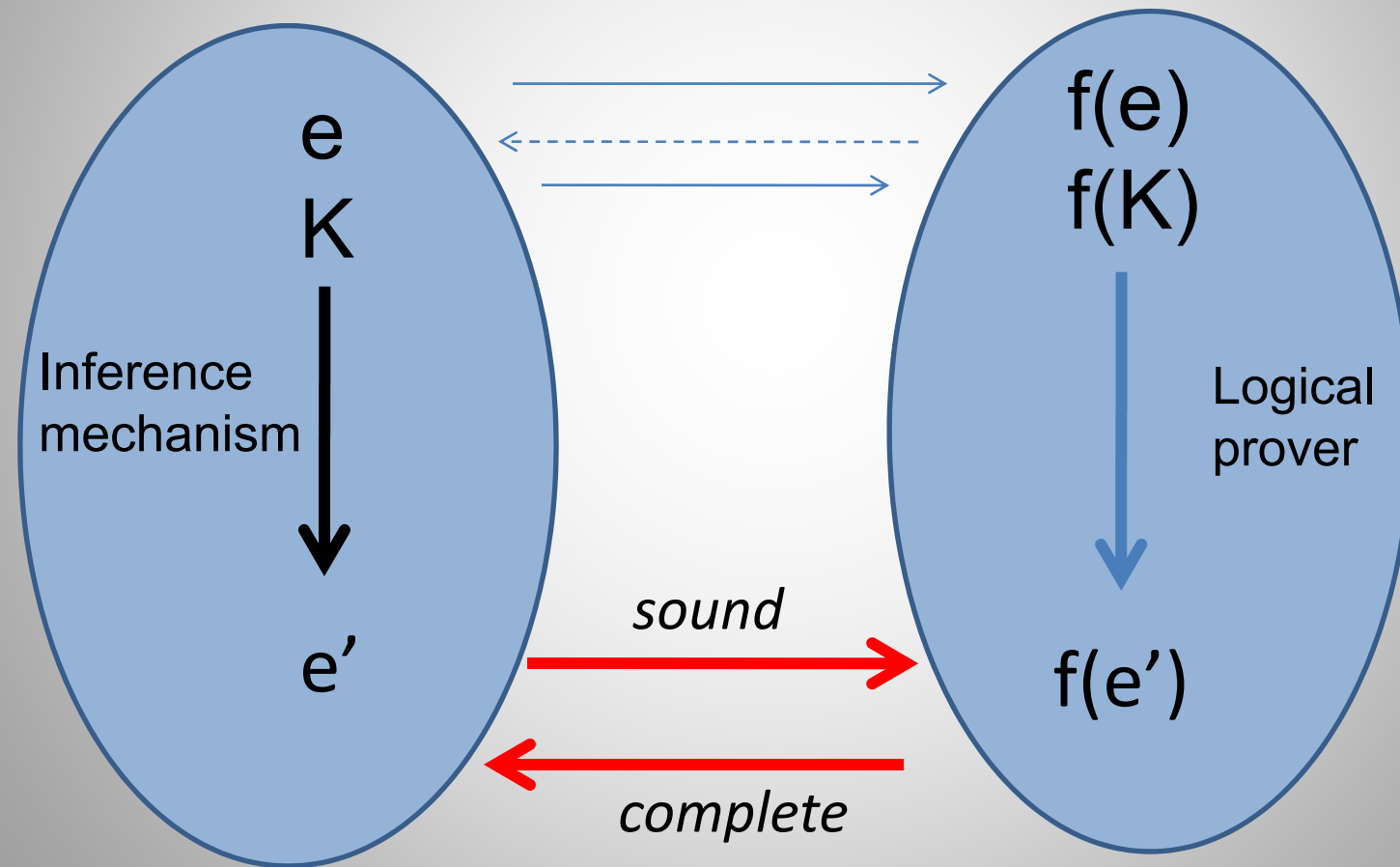
- **Knowledge-based Systems**

- Knowledge Base (Ontology, facts, ...)
- Reasoning mechanism

- **Requirements for a KR&R Formalism**

- Formal semantics (esp. Logical)
- Structured representation of knowledge
- Good computational properties

Knowledge Representation & Reasoning (2)



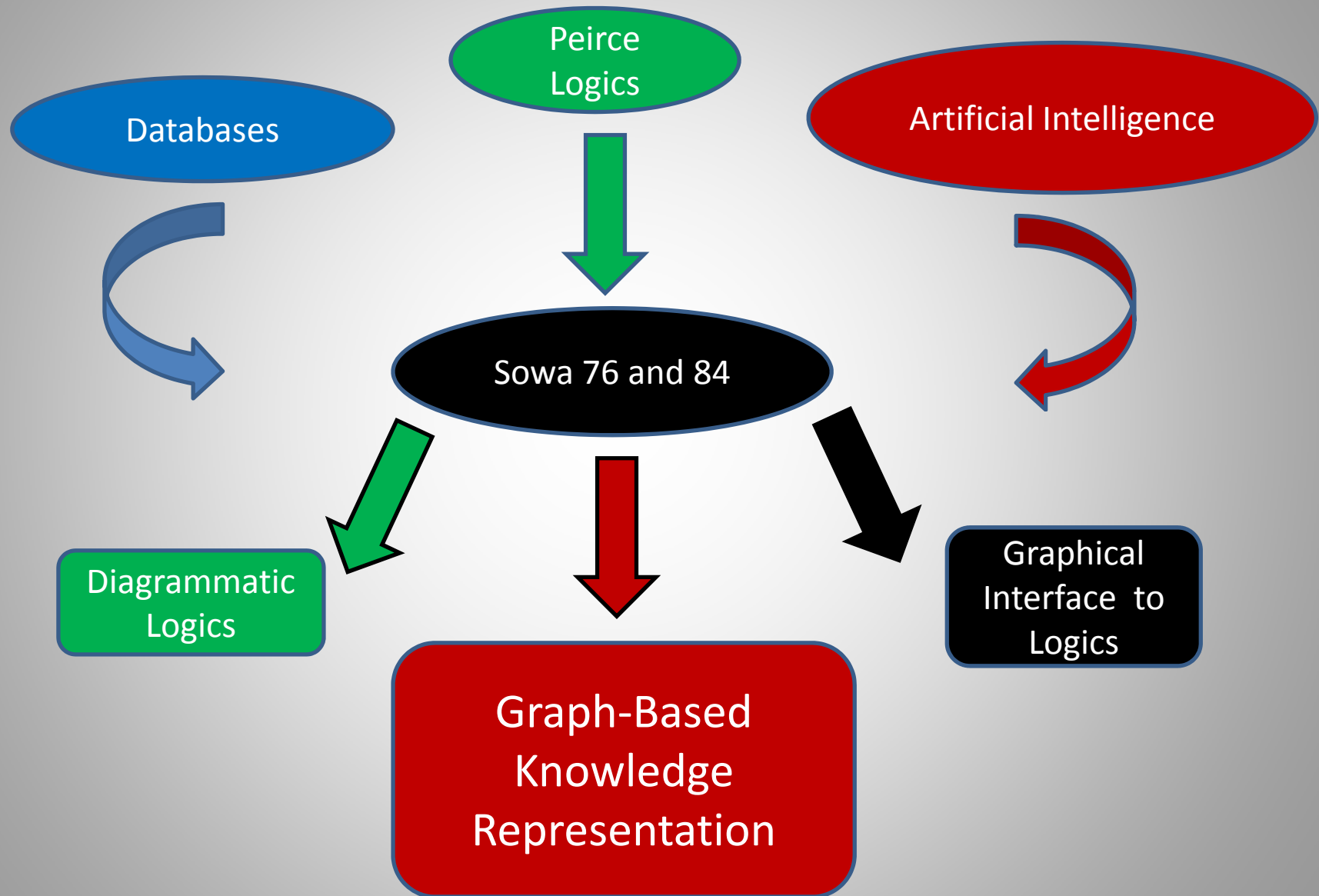
Structured representation of knowledge

- Semantically related pieces of knowledge should be gathered together
- Distinction between ontological knowledge and factual knowledge
- Rules (e.g. for expressing implicit knowledge)
- Constraints

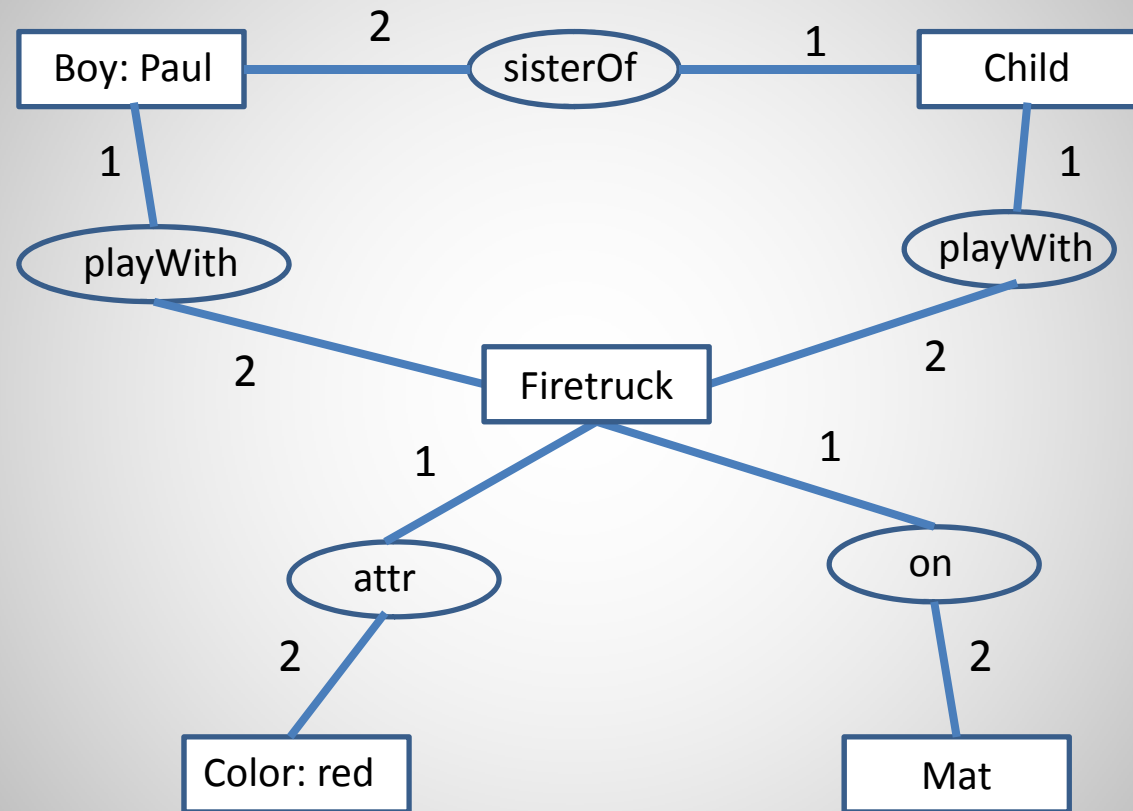
Good computational properties

- Efficient algorithms
- Human computer interface

GBKR Sketchy Genealogy



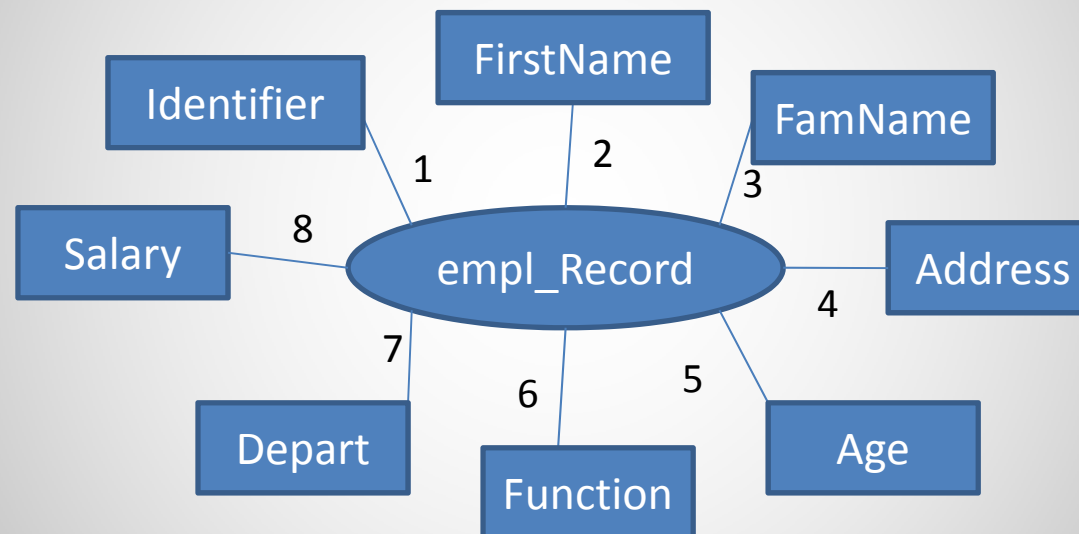
GBKR Basic graph 1



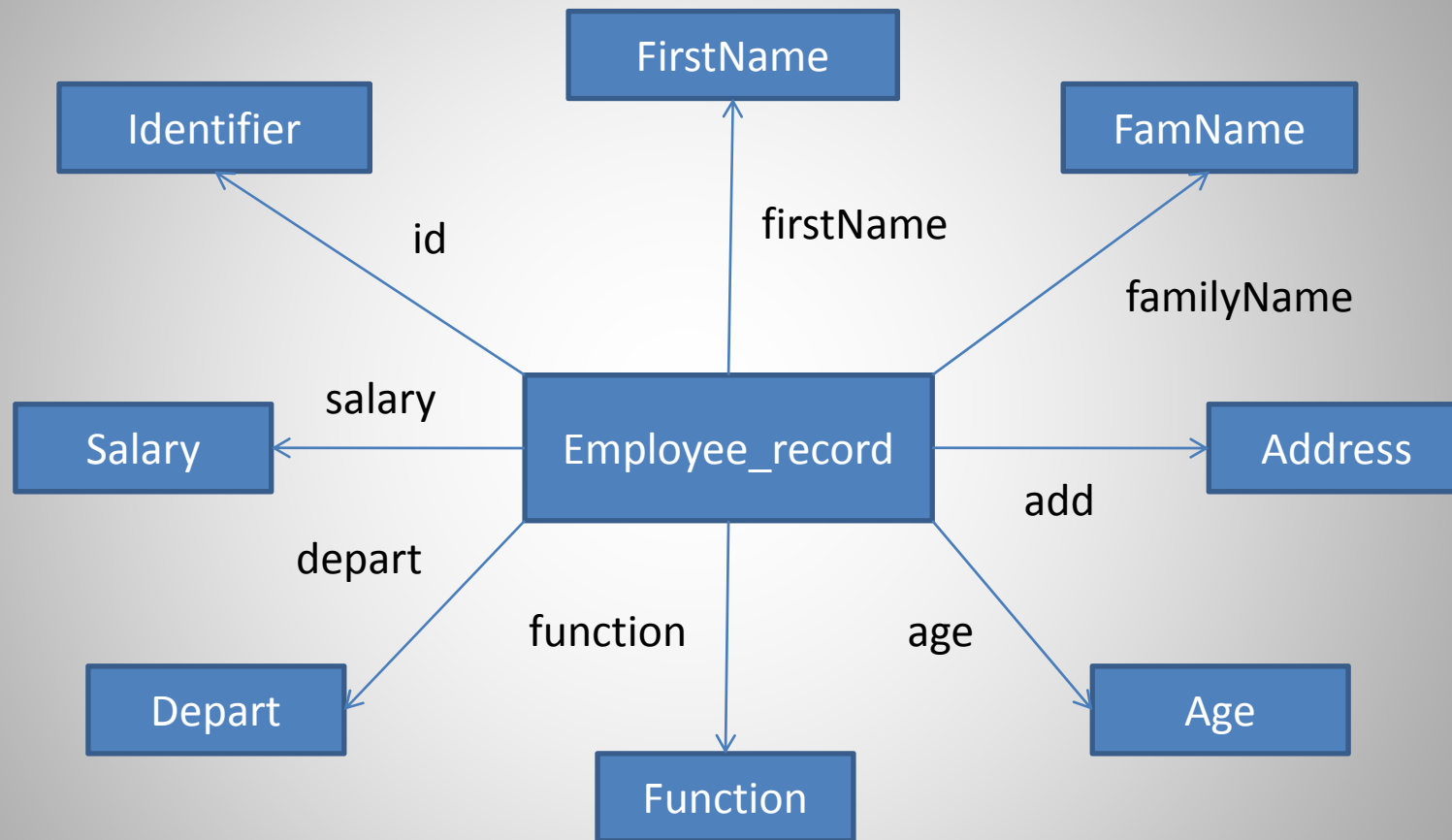
GBKR Basic graph 2



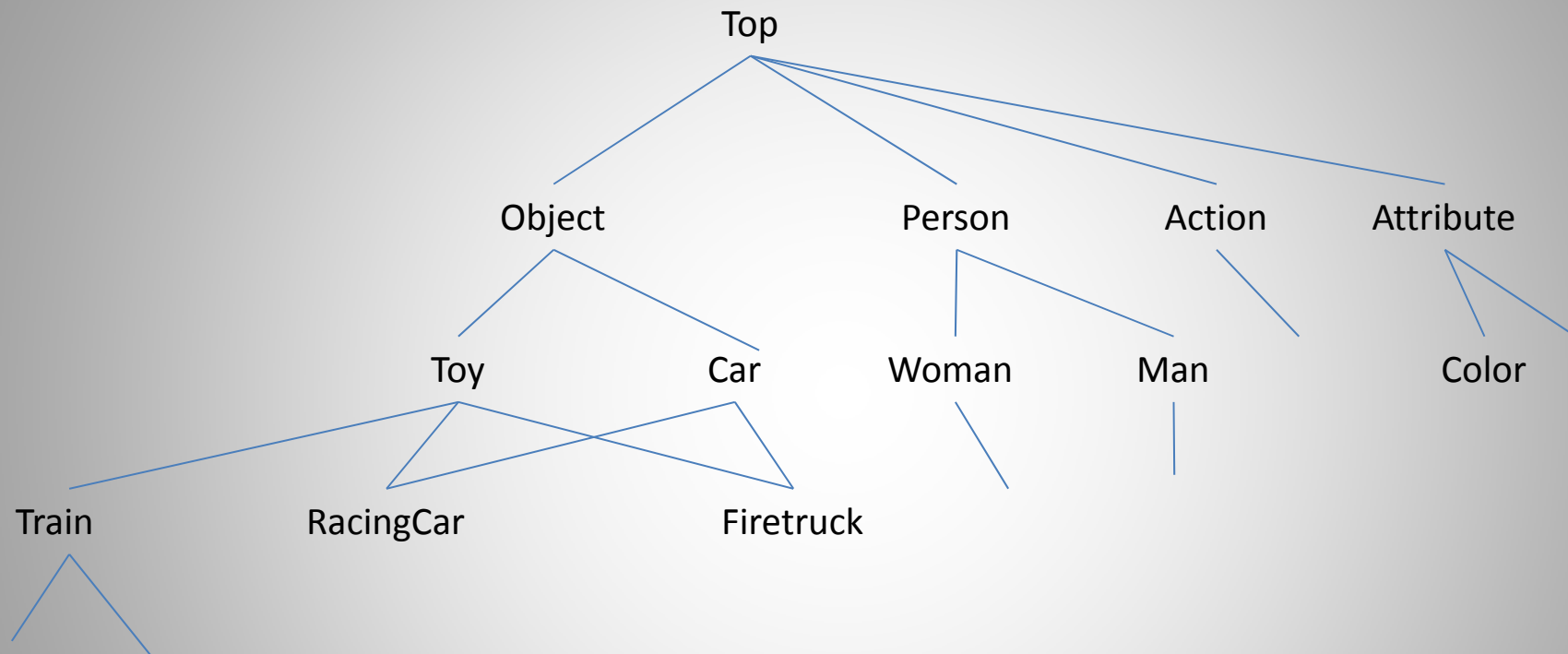
GBKR n-ary relation



GBKR binary role relations

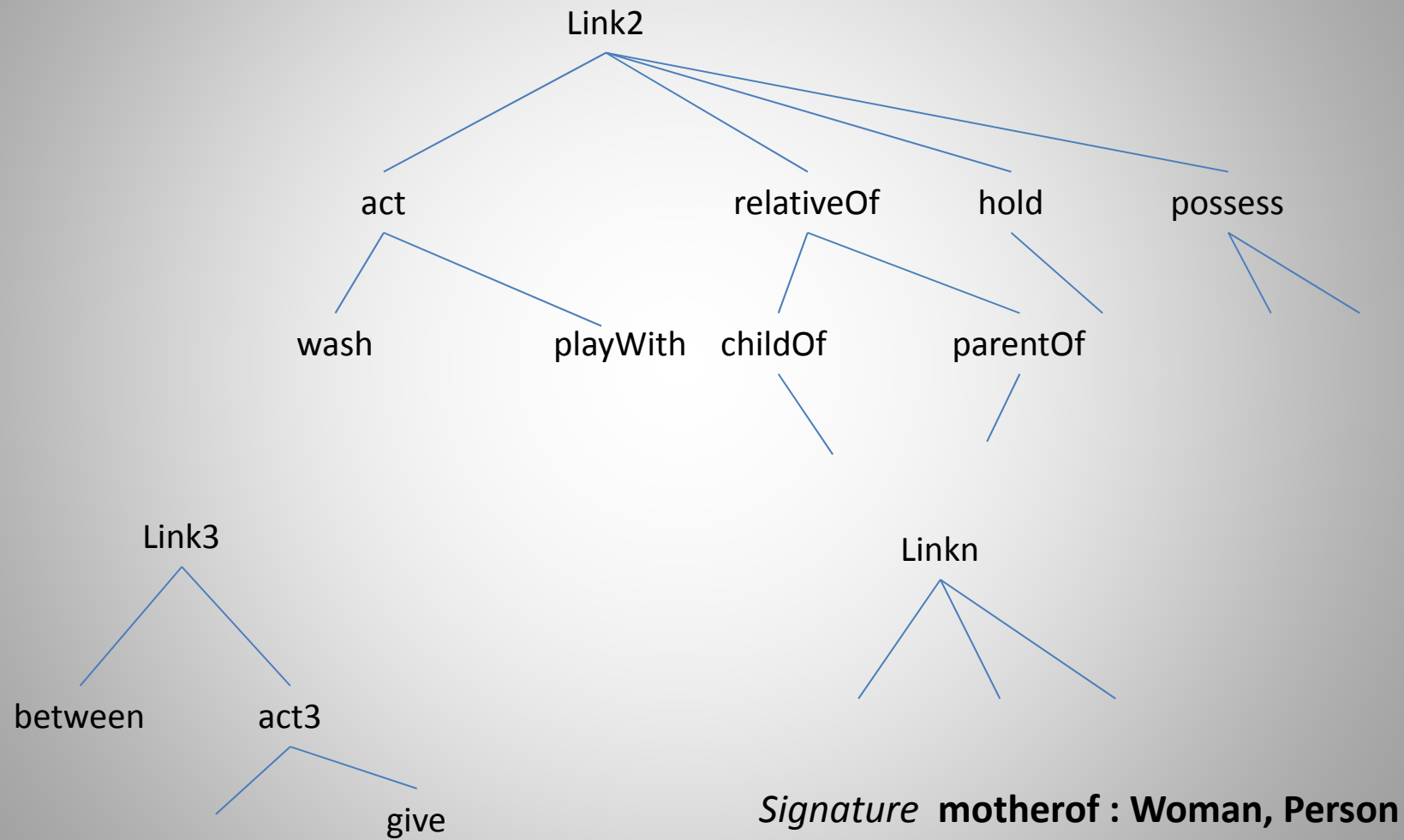


GBKR vocabulary 1



Individuals : Paul is a Boy, Doudou is a Cuddly_Toy, Garfield is a Cat ...

GBKR vocabulary 2

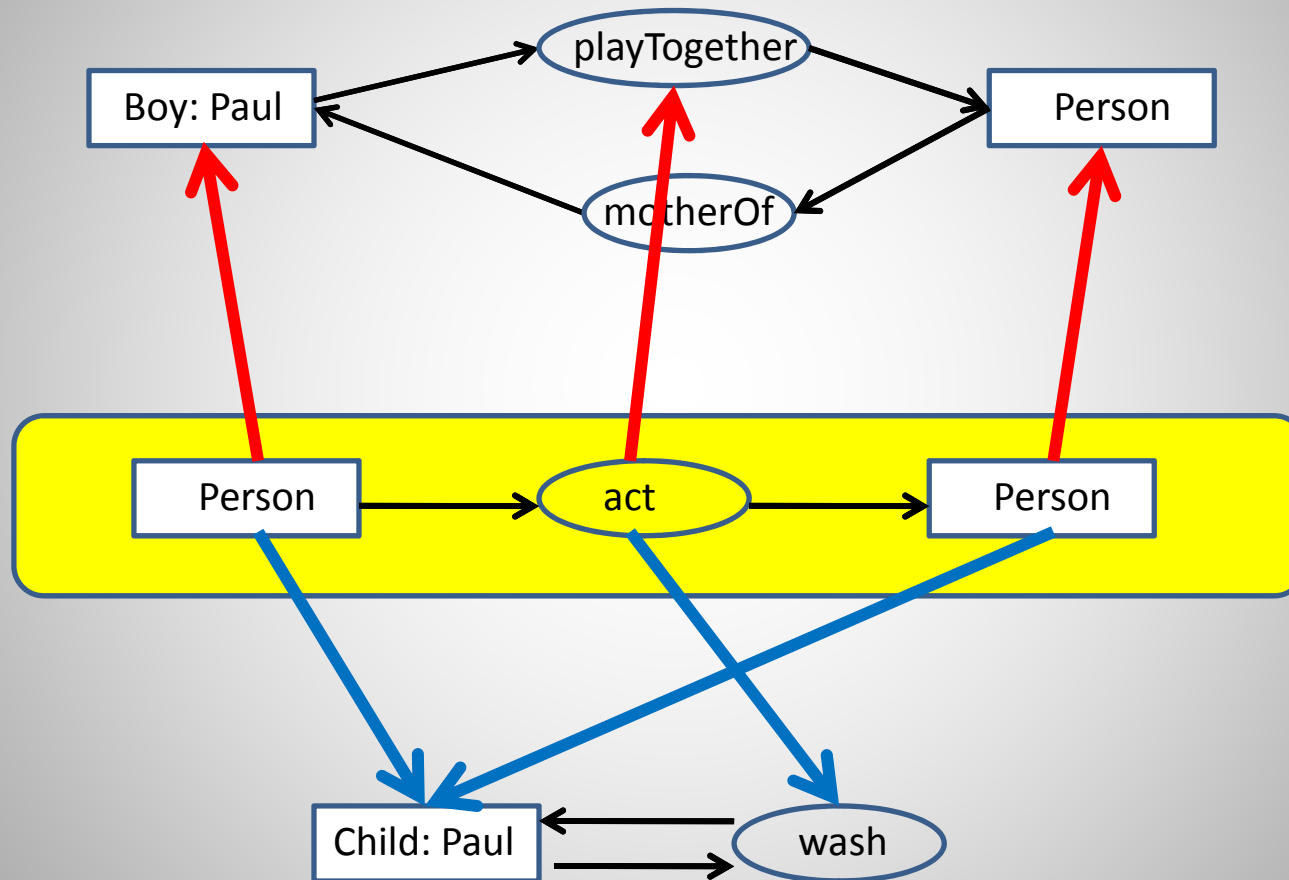


GBKR vocabulary 3

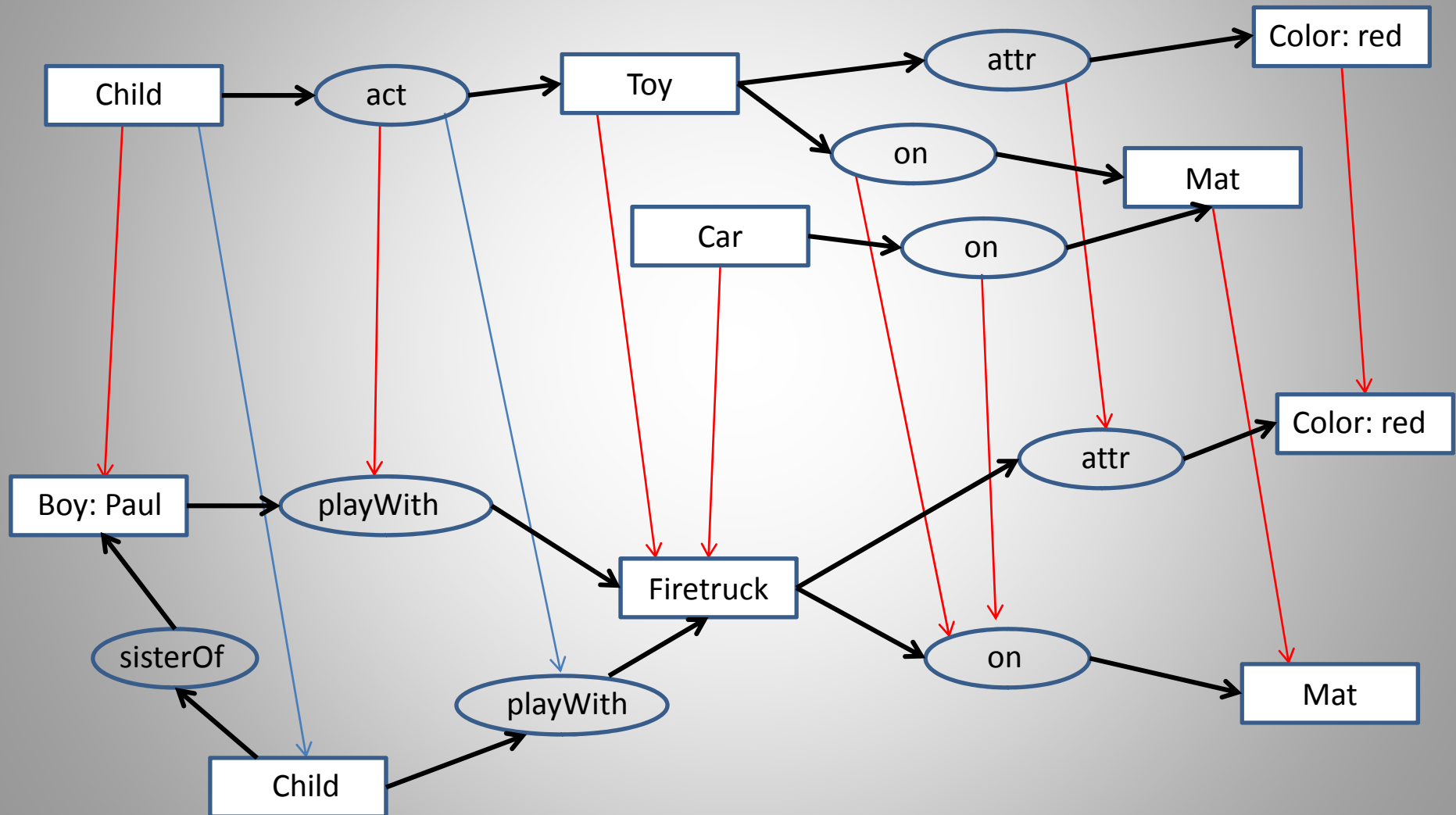
Vocabulary: a **light-weight ontology**

- hierarchies of concept types (classes) and relations
- hierarchical orders: any partial order
- « a kind of » is distinct from « is instance of »
- n-ary relations, $n \geq 1$
- signature of relations (e.g. domain and range for binary)

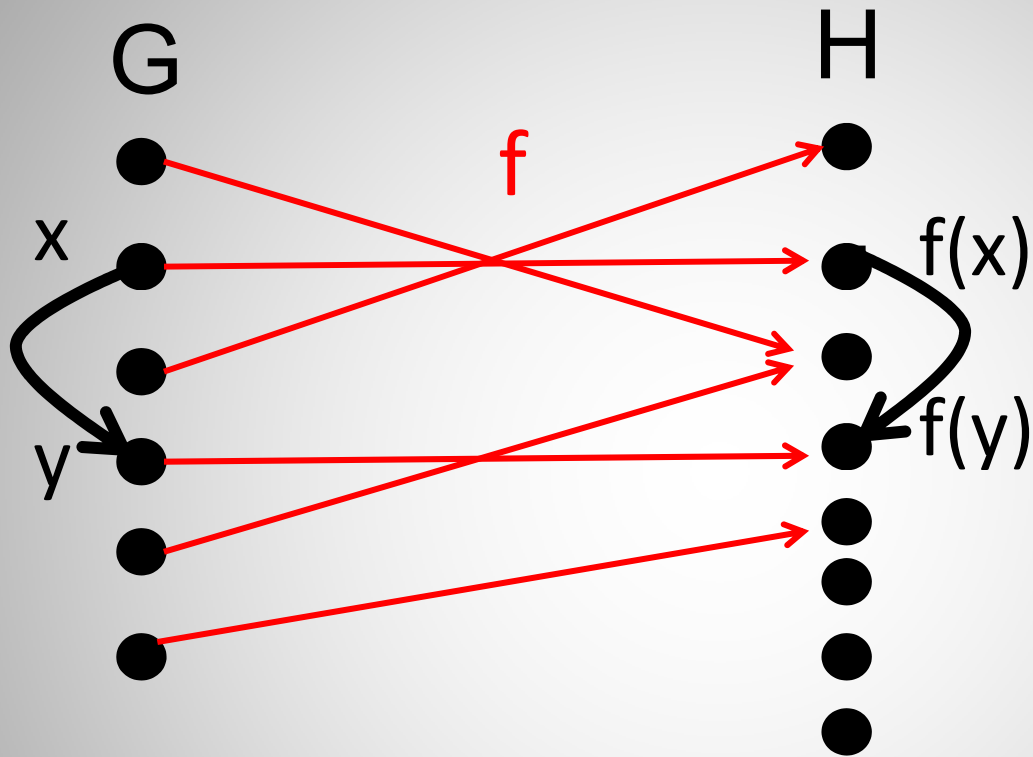
GBKR query/answer 1



GBKR query/answer 2

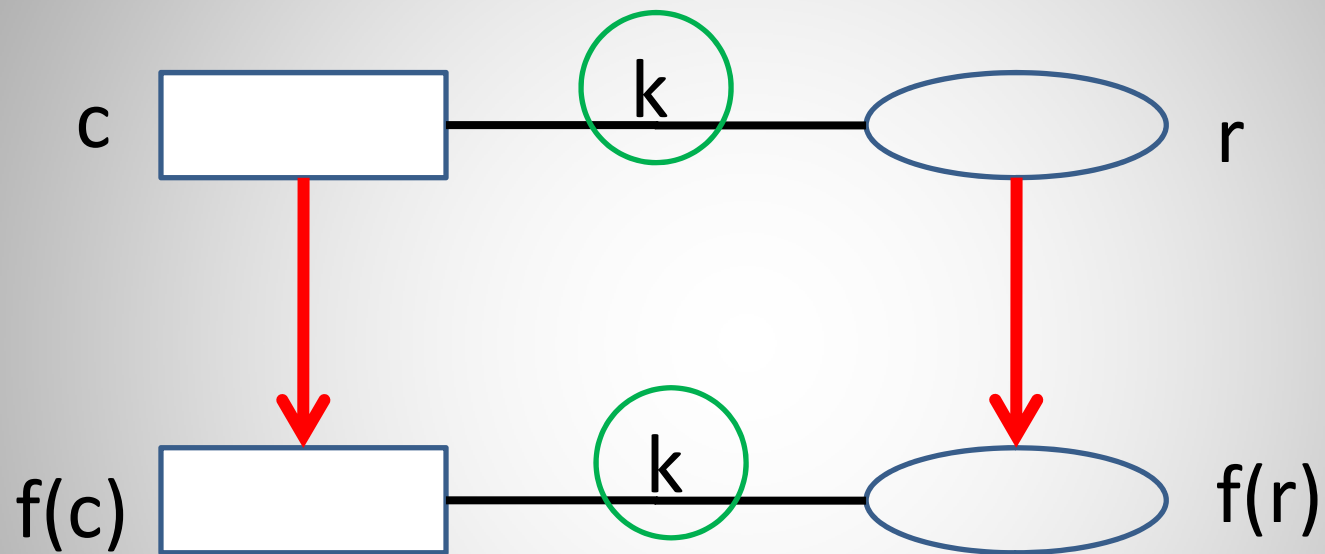


GBKR graph homomorphism 1



xy arc in $G \Rightarrow f(x)f(y)$ arc in H

GBKR graph homomorphism 2



$\text{label}(c) \geq \text{label}(f(c))$ (e.g. $\text{Person} \geq (\text{Boy}, \text{Paul})$)

$\text{label}(r) \geq \text{label}(f(r))$ (e.g. $\text{act} \geq \text{playWith}$)

GBKR graph homomorphism 3

Subsumption relation between graphs

G subsumes H if there is a hom from G to H

Fundamental for structuring a set of graphs

Preorder

Minimal (irredundant) graph

Basis for other operations

Is there an homomorphism from G to H ?

NP-complete

- Efficient backtrack algorithms
- Polynomial cases: G is a tree, tree-width bounded

Polynomial in data complexity

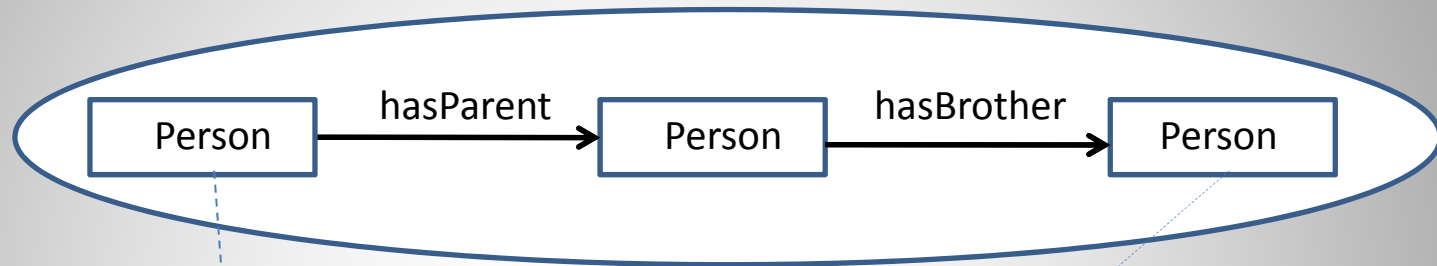
- Size of G not considered

GBKR rules 1

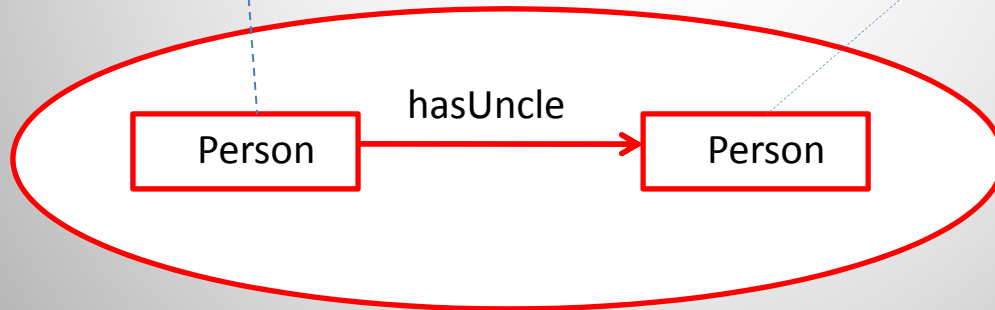
- Fundamental component in knowledge-based systems
- **If** hypothesis H **then** conclusion C
- **If** a piece of knowledge H is present
 then the piece of knowledge C can be added
- Represent general implicit knowledge
- Forward chaining
- Backward chaining

GBKR rules 2

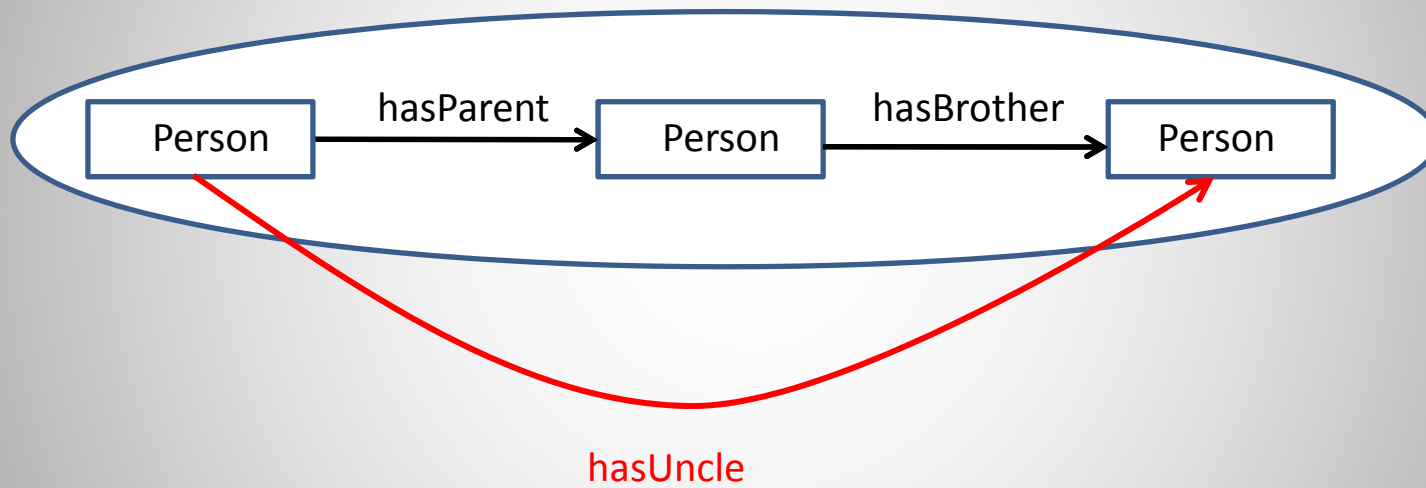
IF



THEN

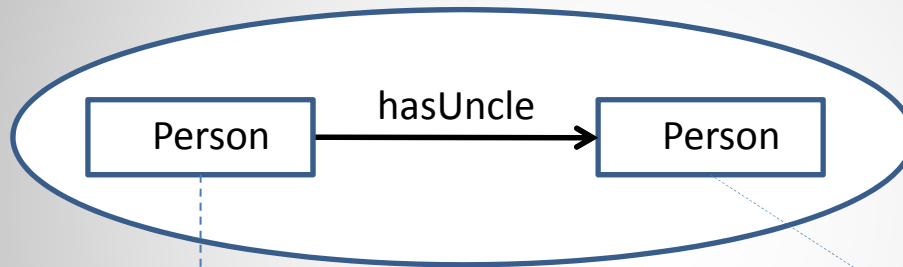


GBKR rules 3

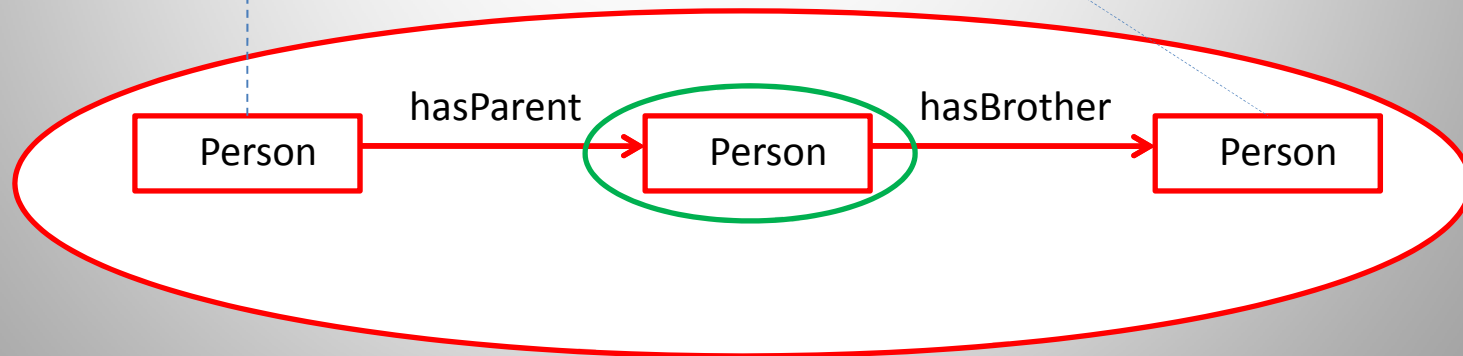


GBKR rules 4

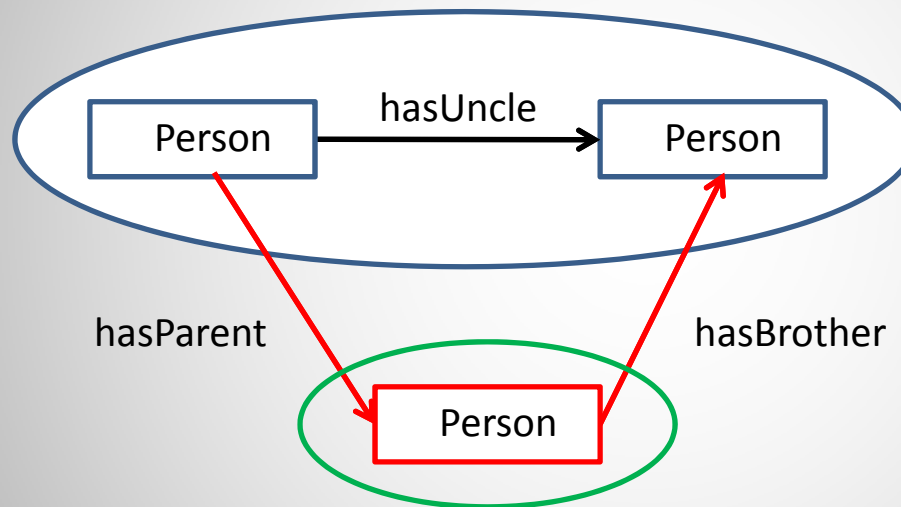
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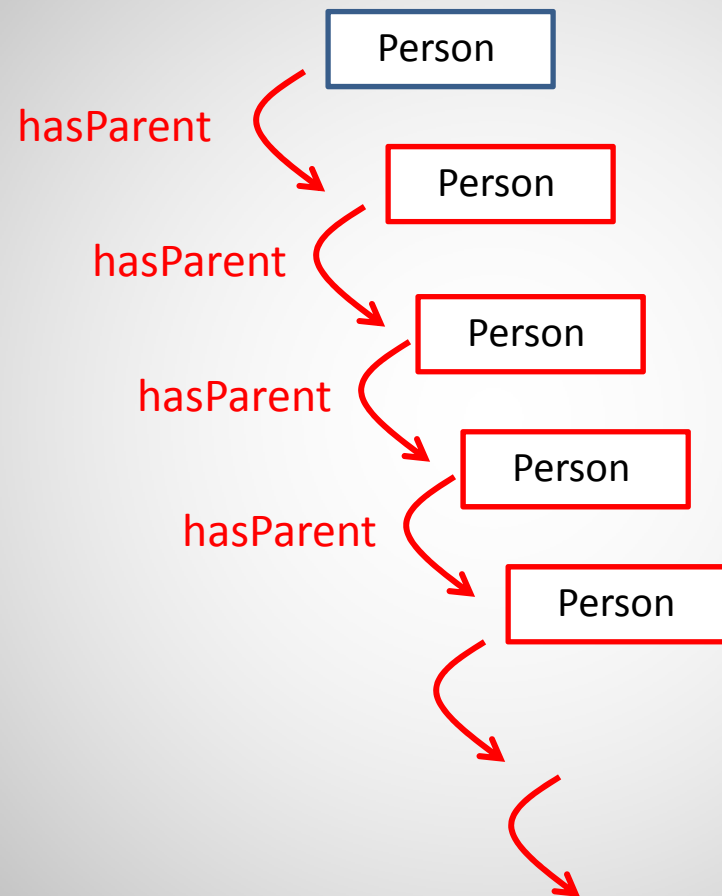
THEN



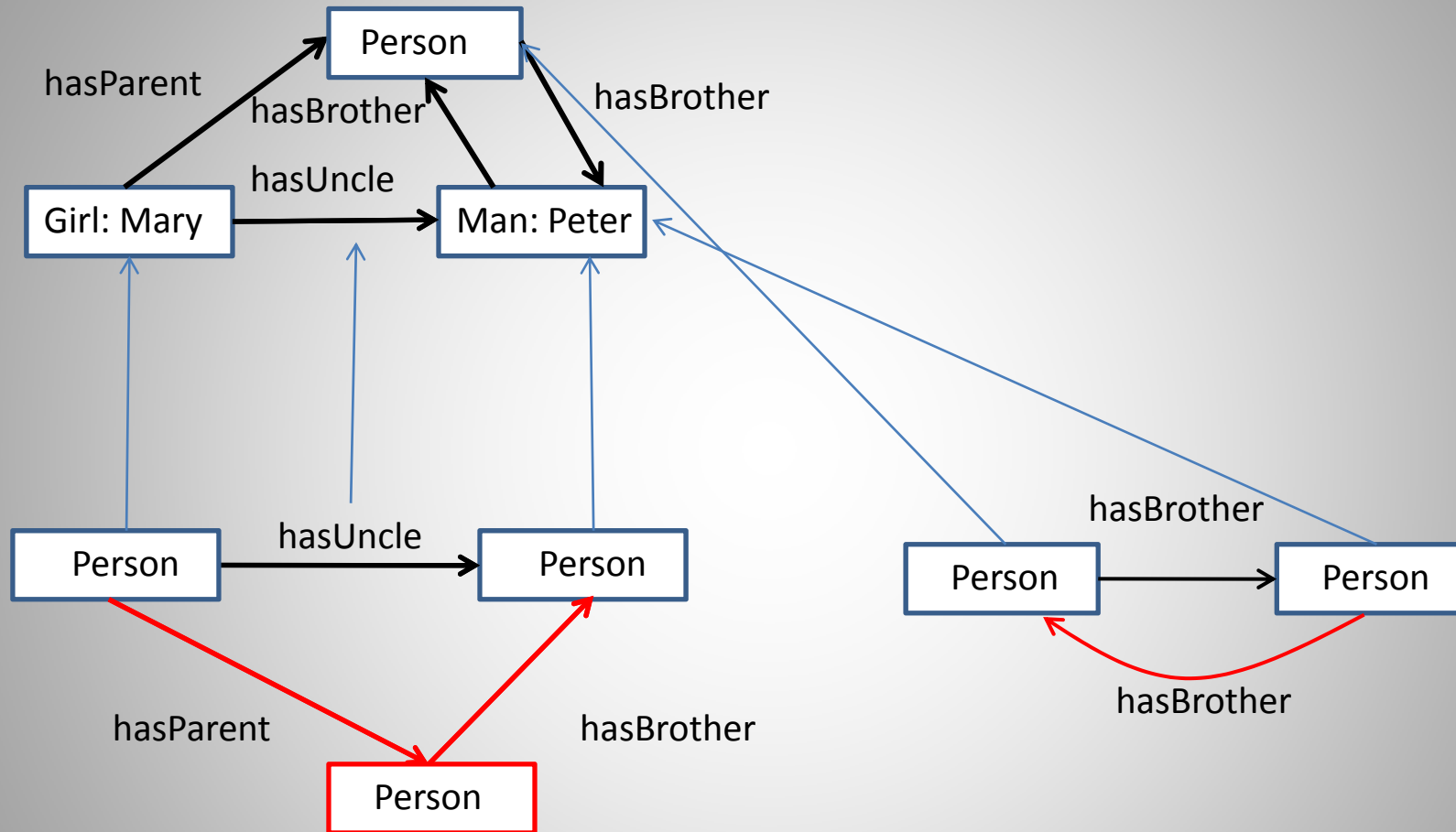
GBKR rules 4



GBKR rules 5



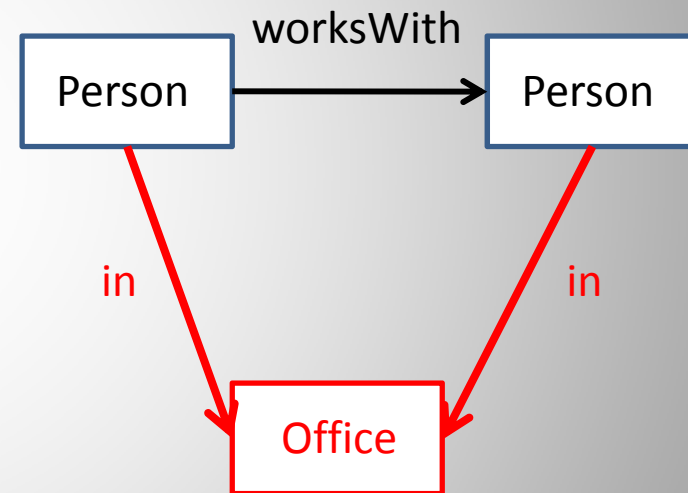
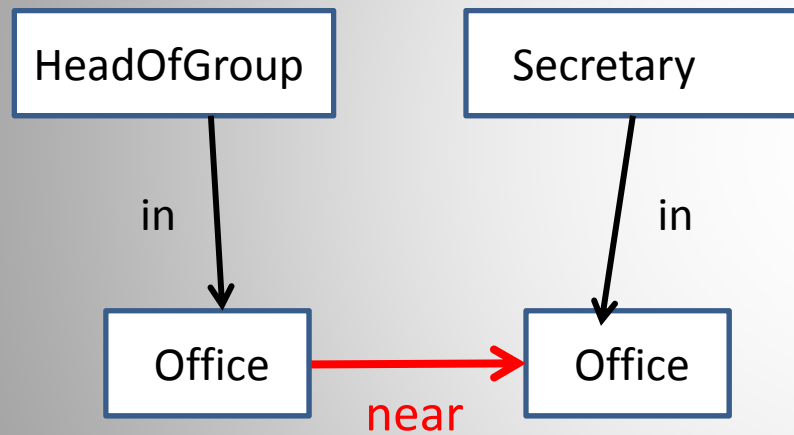
GBKR rules 7



GBKR rules 6

- Equivalent to TGDs in relational databases
- New results obtained with graph viewpoint
- Backward chaining
- Computability/Complexity

GBKR constraints



GBKR other constructs

Conjunctive types

Equality

Nested graphs

Atomic negation

Type definitions

Operations: joins, ...

GBKR hierarchy of models

F an initial world , V a vocabulary, C a set of constraints, R a set of inference rules, E a set of evolution rules (to make evolve a consistent world into a new consistent one)

F is Q deducible from (F,V)?

FC does F satisfies C and is Q deducible from (F,V)?

FR is Q deducible from (F_k, V) ? where F_k is a R-derivation of F

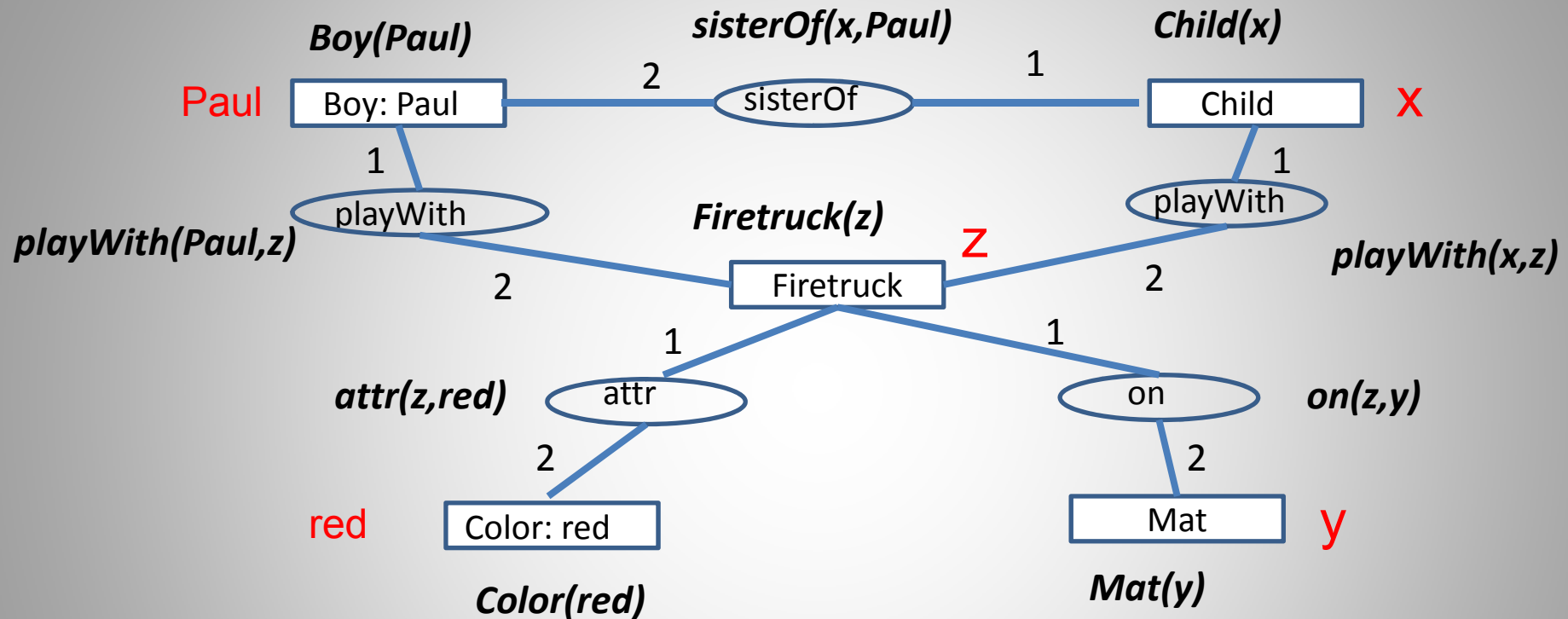
FCR does (F,R) satisfies C and is Q deducible from (F,R)?

FCE does (F,E) satisfies C and is Q deducible from (F,E)?

FRCE deduction pb asks wether F can evolve into a consistent world satisfying the goal Q.

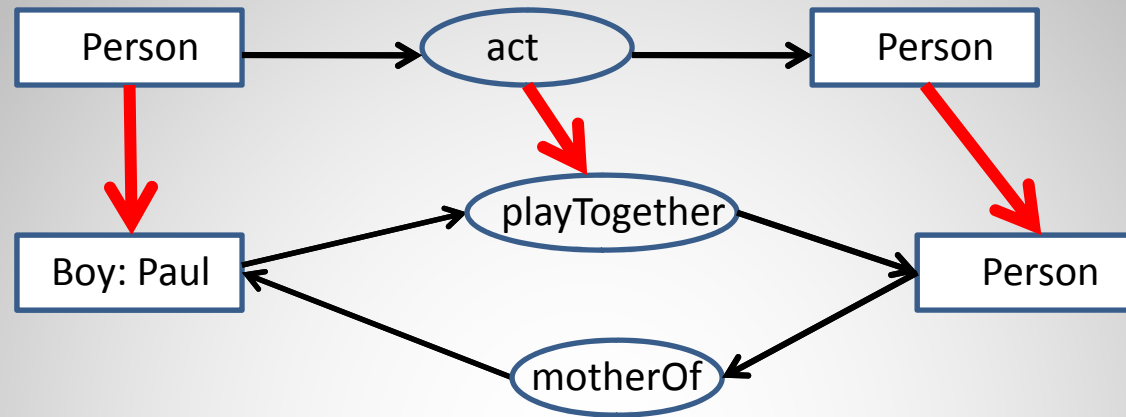
Pause!

GBKR FOL semantics 1

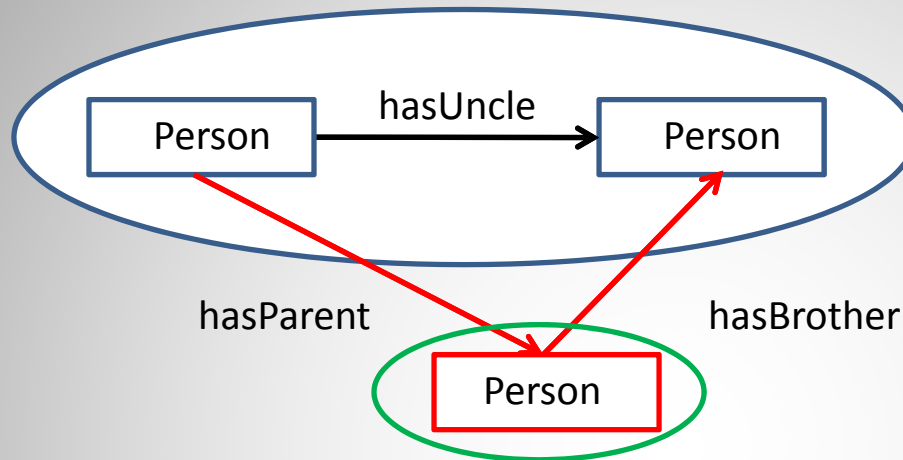


$\Phi(G) = \exists x \exists y \exists z \text{Child}(x) \wedge \text{Mat}(y) \wedge \text{Firetruck}(z) \wedge \text{Boy}(\text{Paul}) \wedge \text{Color}(\text{red}) \wedge \text{sisterOf}(x, \text{Paul}) \wedge \text{playWith}(\text{Paul}, z) \wedge \text{playWith}(x, z) \wedge \text{attr}(z, \text{red}) \wedge \text{on}(z, y)$

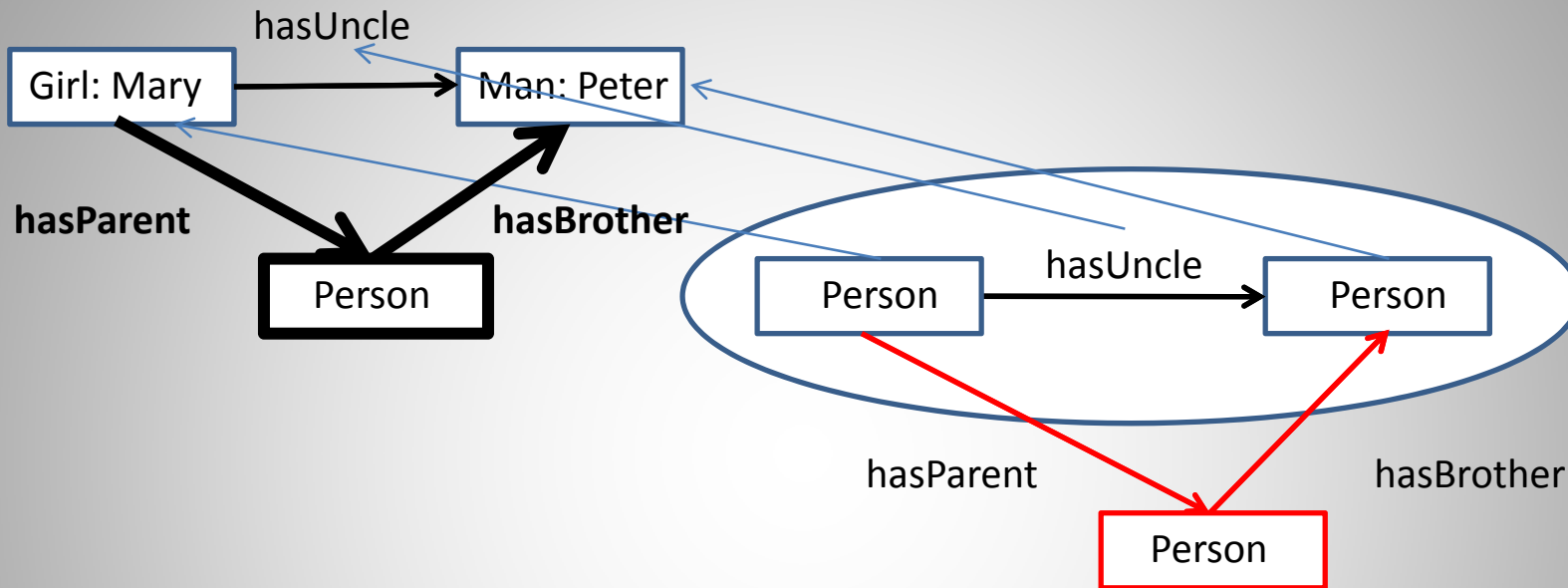
GBKR FOL semantics 2


$$(\forall x \text{Boy}(x) \rightarrow \text{Person}(x)) \wedge (\forall x \forall y \text{playTogether}(x,y) \rightarrow \text{act}(x,y)) \wedge$$
$$\exists x (\text{Boy}(\text{Paul}) \wedge \text{Person}(x) \wedge \text{playTogether}(\text{Paul}, x) \wedge \text{motherOf}(x, \text{Paul}))$$
$$\rightarrow$$
$$\exists x \exists y \text{Person}(x) \wedge \text{Person}(y) \wedge \text{act}(x,y)$$

GBKR FOL semantics 3


$$\forall x \forall y \left((\text{Person}(x) \wedge \text{Person}(y) \wedge \text{hasUncle}(x,y)) \rightarrow \right. \\ \left. \exists z \text{Person}(z) \wedge \text{hasParent}(x,z) \wedge \text{hasBrother}(z,y) \right)$$

GBKR FOL semantics 4



$$\forall x(\text{Girl}(x) \rightarrow \text{Person}(x)) \wedge \forall x(\text{Man}(x) \rightarrow \text{Person}(x)) \wedge$$

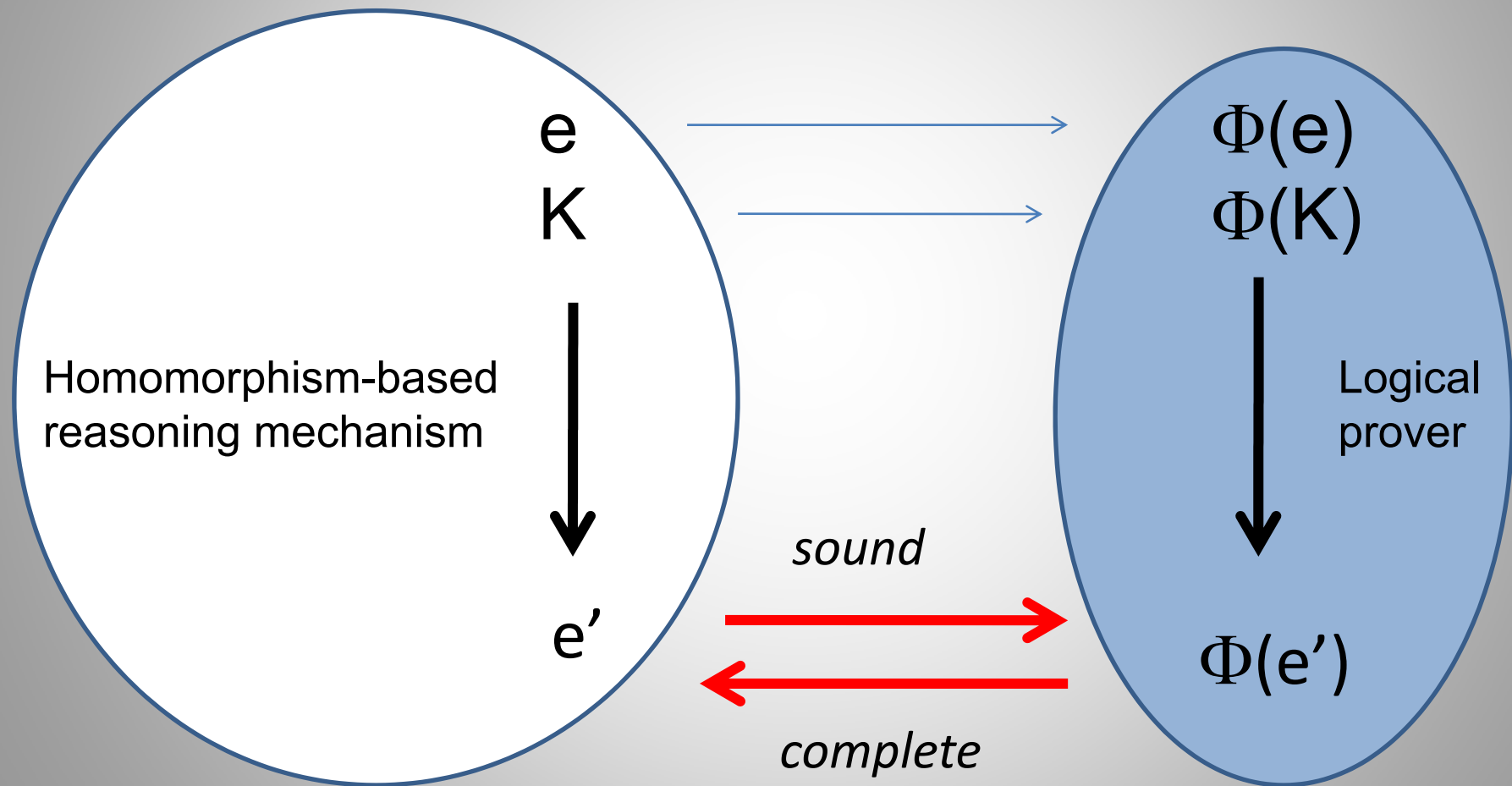
$$\forall x \forall y (\text{Person}(x) \wedge \text{Person}(y) \wedge \text{hasUncle}(x, y) \rightarrow \exists z \text{Person}(z) \wedge \text{hasParent}(x, z) \wedge \text{hasBrother}(z, y)) \wedge$$

$$\text{Girl}(\text{Mary}) \wedge \text{Man}(\text{Peter}) \wedge \text{hasUncle}(\text{Mary}, \text{Peter})$$

→

$$\exists z \text{Person}(z) \wedge \text{hasParent}(\text{Mary}, z) \wedge \text{hasBrother}(z, \text{Peter})$$

GBKR Soundness and completeness



GBKR equivalent formalisms 2

Large variety of graphs

Directed graphs, undirected graphs,
multigraphs, labeled graphs, hypergraphs

Transformations of one kind to another

Polynomial reductions for homomorphism problems

Parsimonious equivalence between some of them especially:

Homomorphism for GBKR graphs and

Homomorphism for simple unlabeled (directed or undirected) graphs

Relational structures and relational databases

Query evaluation problem

Instance: a database instance D and a conjunctive query q

Question: Does D contain an answer to q ?

Query containment problem

Instance: two queries q and q' ?

Question: for any D does $q(D)$ contain $q'(D)$?

These pbs are polynomially equivalent to GBKR graph hom

GBKR equivalent formalisms 4

Constraint satisfaction problem

Variables x_1, \dots, x_n

Domains D_1, \dots, D_n

Constraints C_1, \dots, C_k

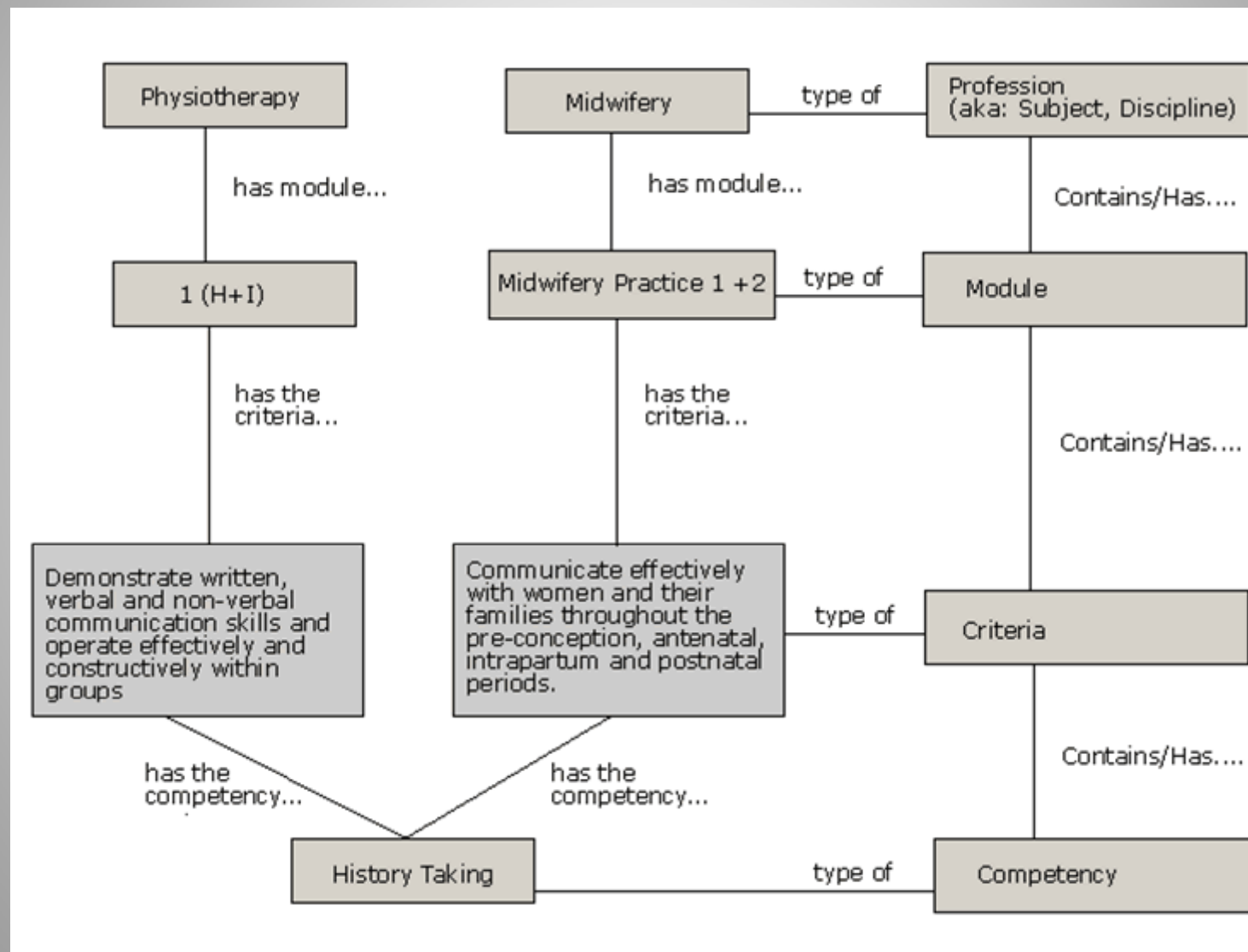
Solution: an assignment of the variables satisfying the constraints

Parsimonious reduction between CSP and BH-homomorphism

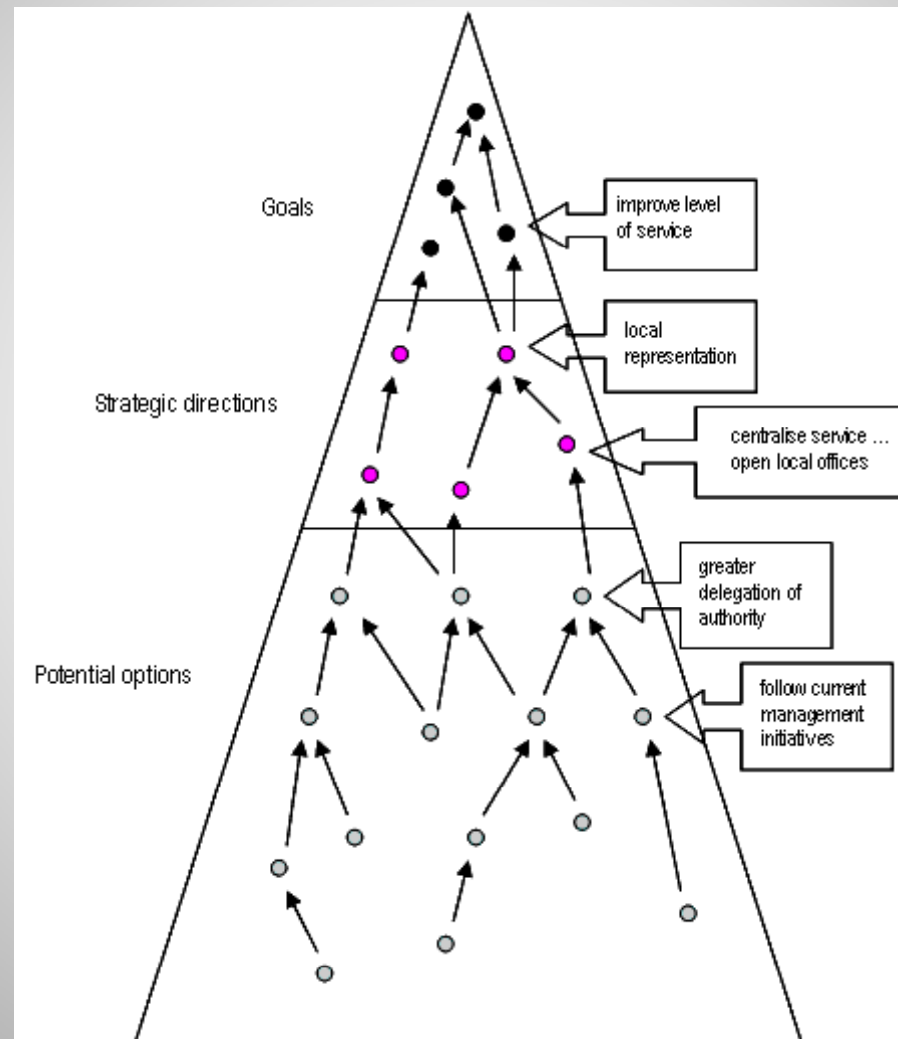
Graph(ical) Models

- Semantic Networks
- Entity-Relationship Model
- KL-One
- UML

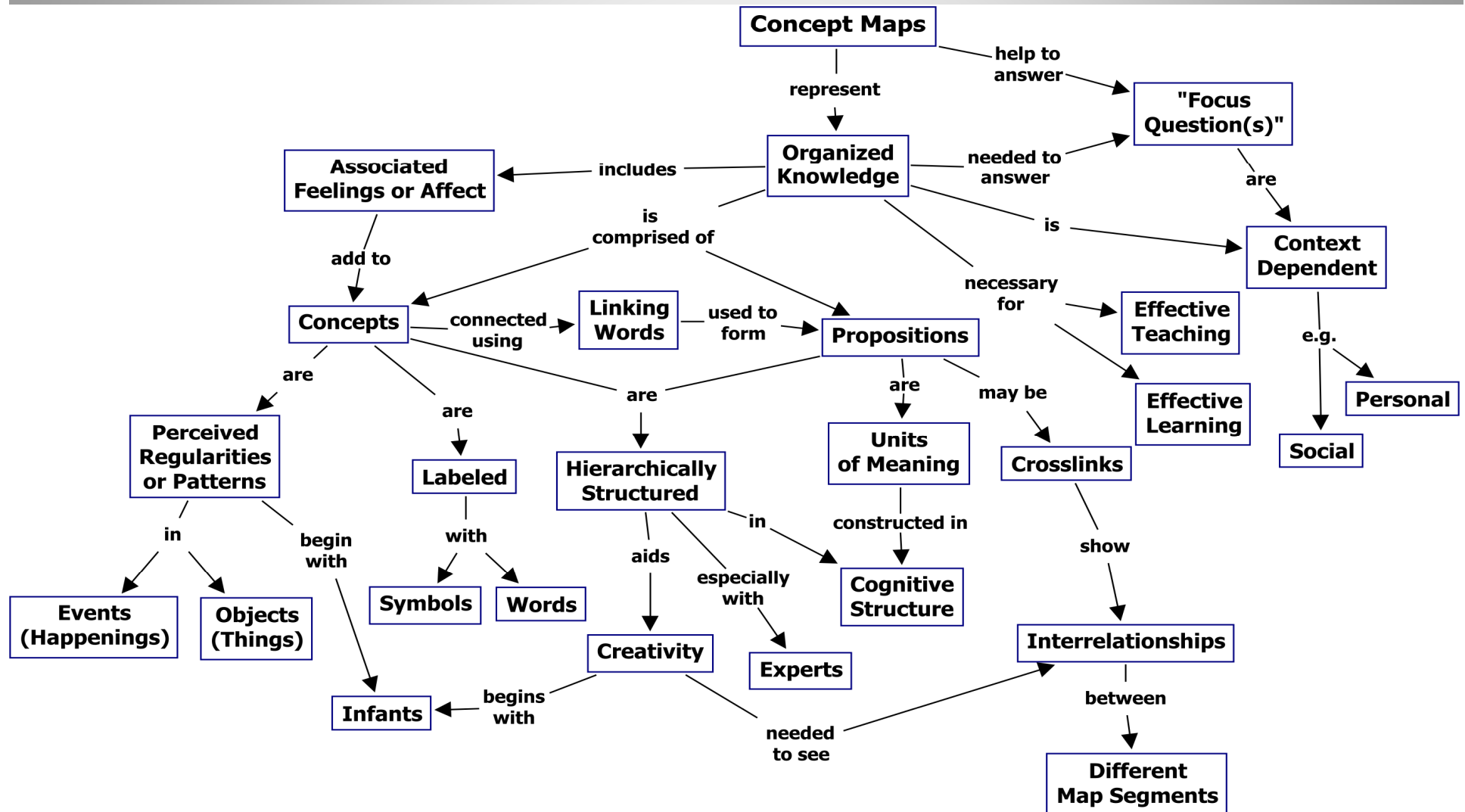
Topic Maps



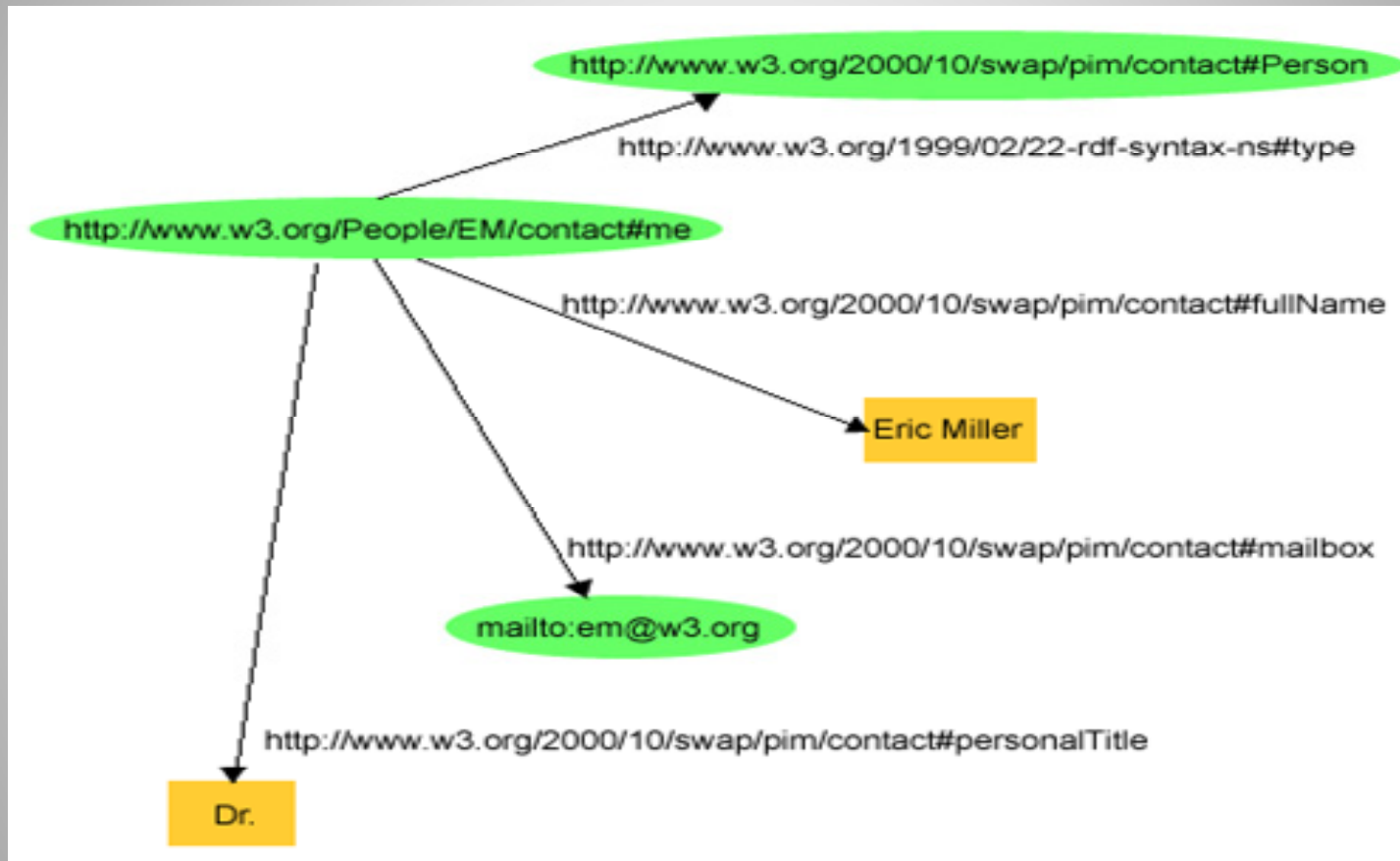
Cognitive Maps



Concept Maps



RDF/S



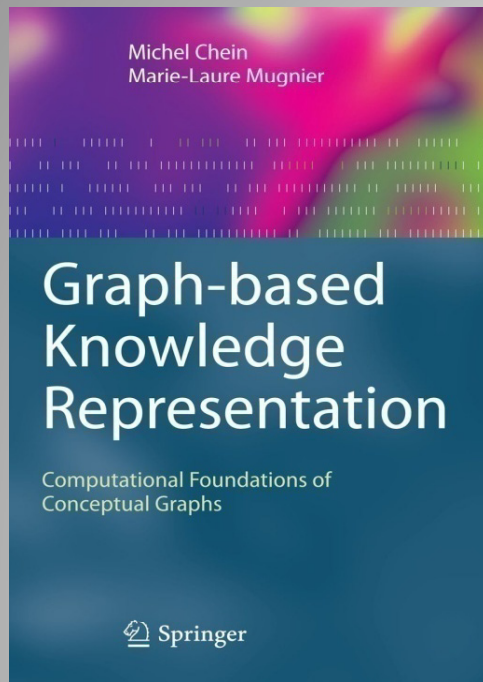
Strengths

- Formal semantics (FOL or model theory)
- Structured representation of knowledge
- Good computational properties
- Numerous algos
- Diagrammatic/Linear
- Visual aspects GUI easily interpreted
- Support for intuition

Weaknesses

- Diagrammatic/Linear
 - Ambiguity in interpretation, a formal (linear !) semantics is needed
- Visual aspects
 - difficulty large or automatically built graphs
- Support for intuition
 - imprecision (modality)

Difficulty to build a relevant ontology



COGUI <http://www.lirmm.fr/cogui/>

GUI, developed in Java, for building GBKR knowledge bases

COGXML format, COGITANT compatible

There is a translator from and to RDF(S)

COGITANT <http://cogitant.sourceforge.net/>

Library of C++ classes build applications based on the GBKR model.
Classes for each object of the model (vocabulary, graph, rule, constraint...)

and for the main operations of the model (homomorphism, application of rules...).

Thank you for your attention!

GBKR equivalent formalisms 1

Parsimonious reduction

A polynomial reduction from $P1$ to $P2$ is ***parsimonious*** if the number of solutions of any YES instance of $P1$ is equal to the number of solutions of its corresponding YES instance of $P2$

Parsimonious equivalence

Two problems $P1$ and $P2$ are ***parsimoniously equivalent*** if each of them is parsimoniously reducible to the other

Description Logics

- Rooted in frames and semantic networks
- Remedy critiques on their ancestors
 - distinction between ontology and facts
 - provided with FOL semantics
- Small intersection
 - rooted trees with binary relations
 - DL tailored for the comparison ELIRO1

Difficulty to build a knowledge base manually and a fortiori automatically

GBKR hierarchy of models

- \mathcal{F} is Q deducible from $(\mathcal{F}, \mathcal{V})$?
- \mathcal{F} describes an initial world and \mathcal{V} a vocabulary
- \mathcal{FC} does \mathcal{F} satisfies C and is Q deducible from $(\mathcal{F}, \mathcal{V})$?
- \mathcal{C} constraints define the consistency of a world
- \mathcal{FR} is Q deducible from $(\mathcal{F}_k, \mathcal{V})$ where \mathcal{F}_k is a \mathcal{R} -derivation of \mathcal{F} ?
- \mathcal{R} inference rules complete the description of any world
- \mathcal{FCR} does $(\mathcal{F}, \mathcal{R})$ satisfies C and is Q deducible from $(\mathcal{F}, \mathcal{R})$?
- \mathcal{FCE} does $(\mathcal{F}, \mathcal{E})$ satisfies C and is Q deducible from $(\mathcal{F}, \mathcal{E})$?
- \mathcal{E} evolution rules try to make evolve a consistent world into a new consistent one

\mathcal{FRCE}

The deduction pb asks wether \mathcal{F} can evolve into a consistent world satisfying the goal Q.